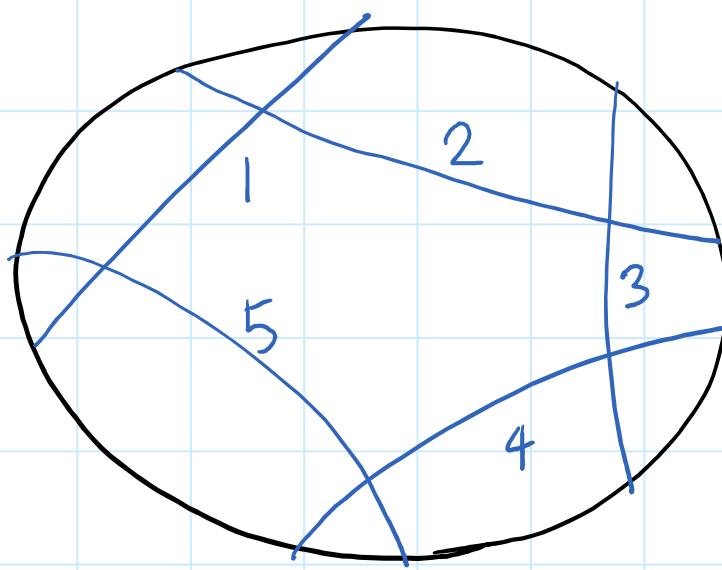


Recognizing Circle Graphs

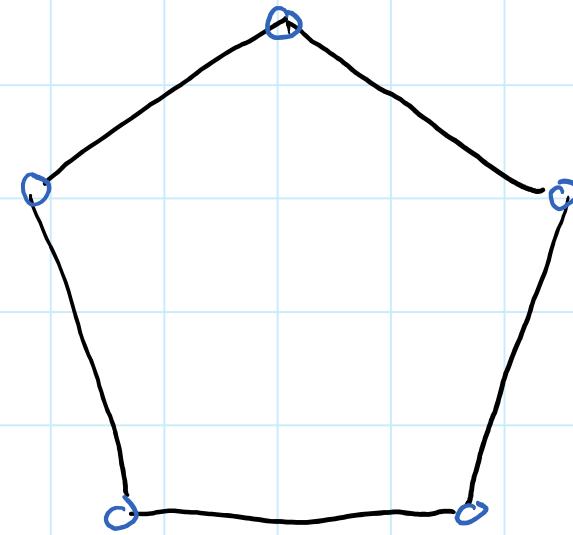
Jim Geelen

Edward Lee

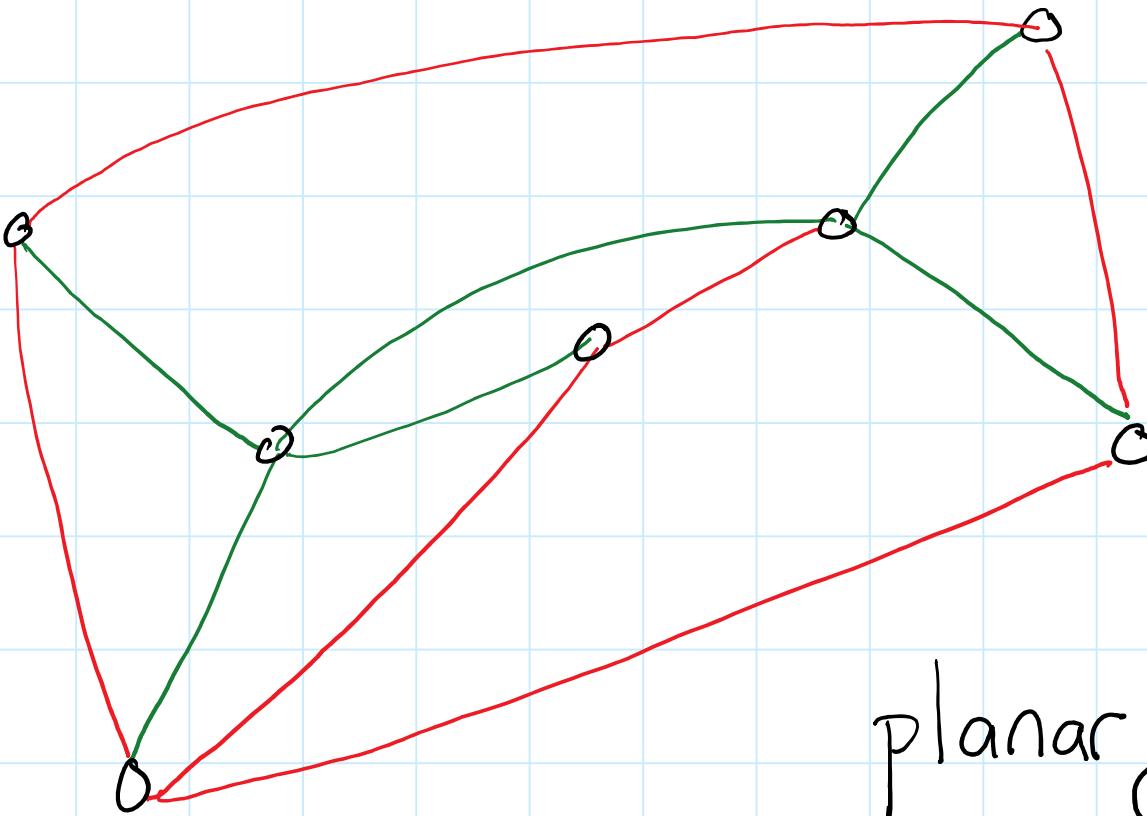
University of Waterloo



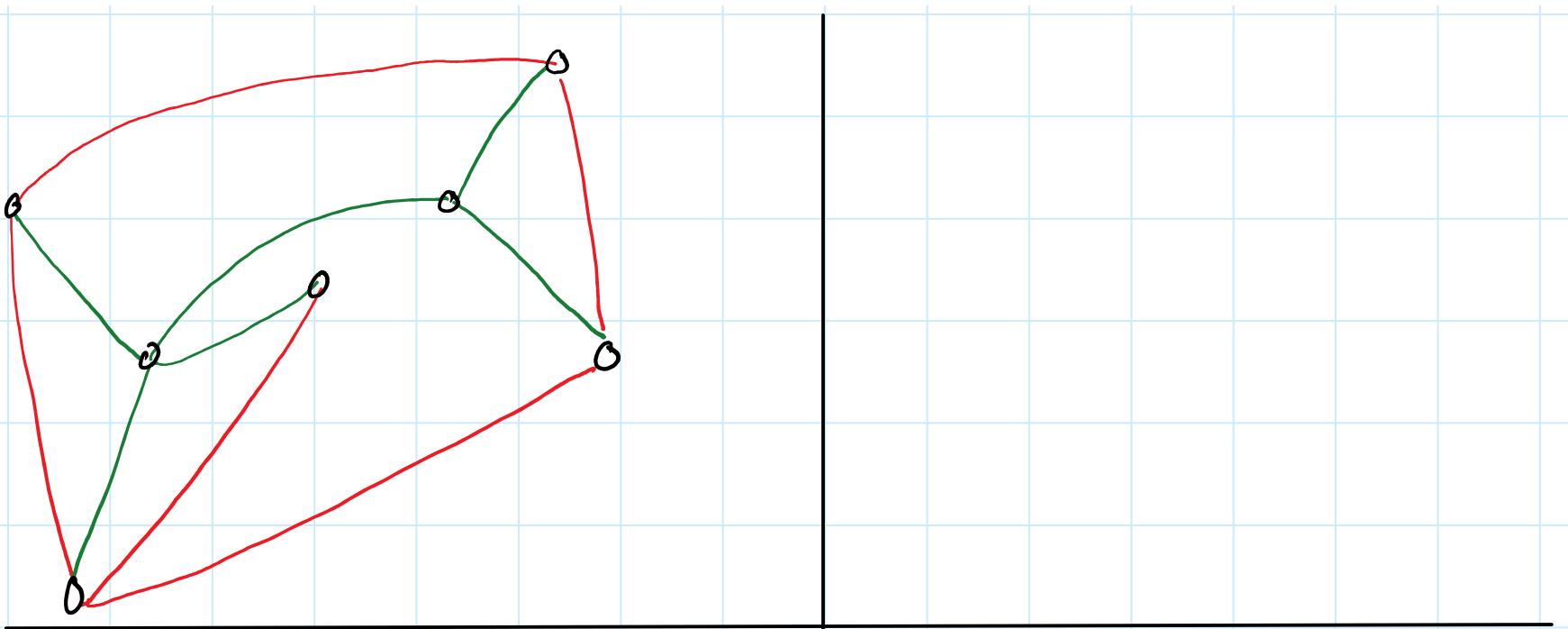
Chord Diagram
C

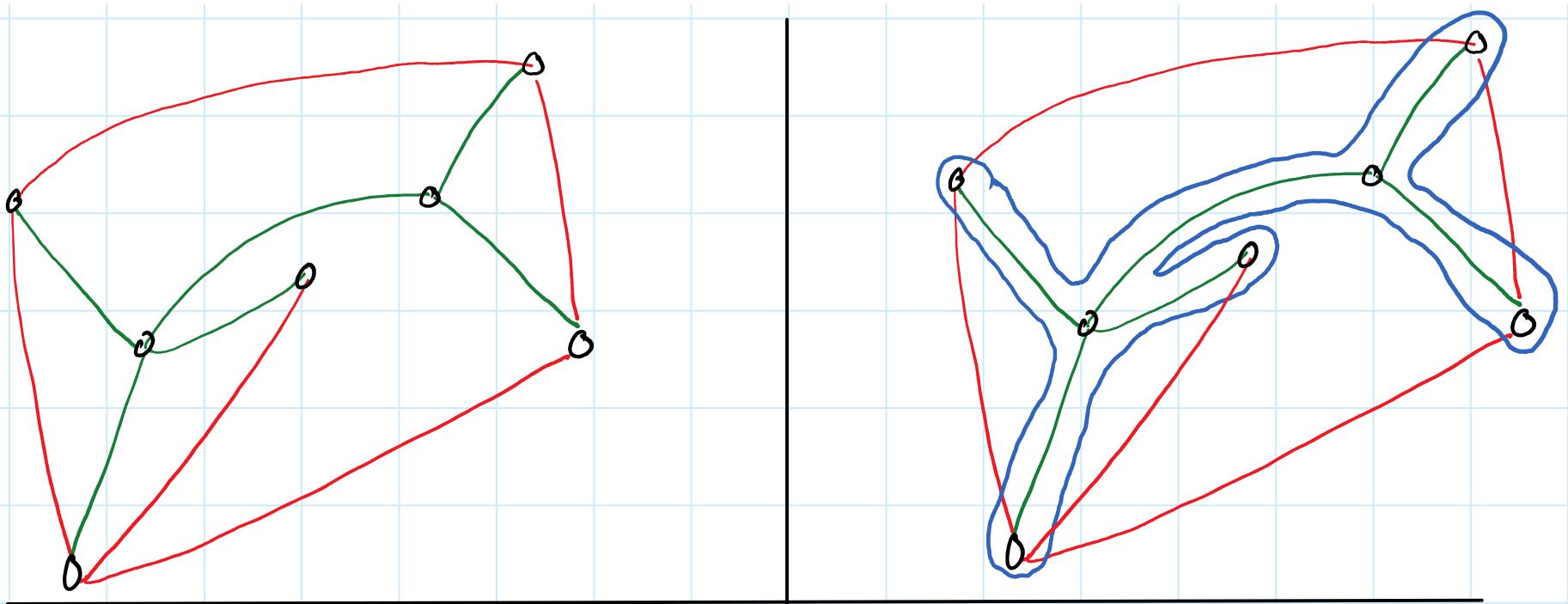


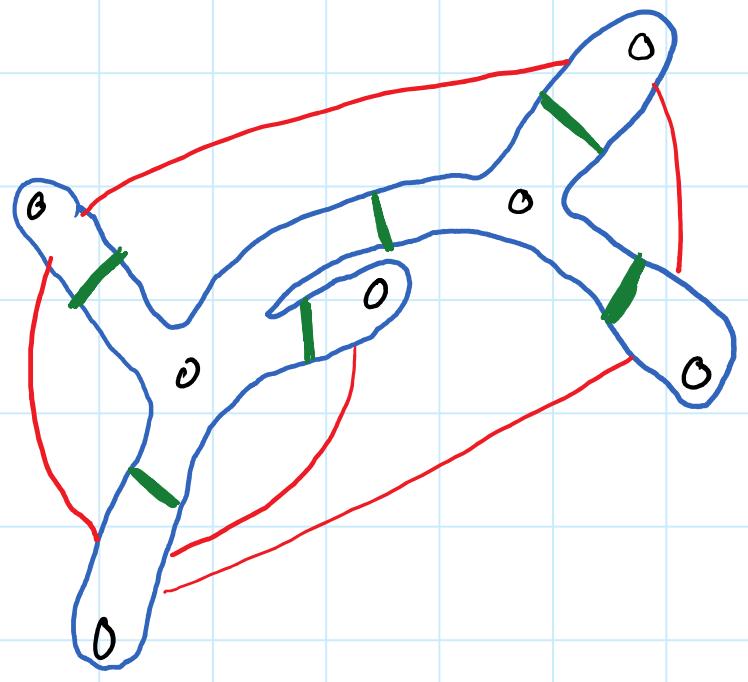
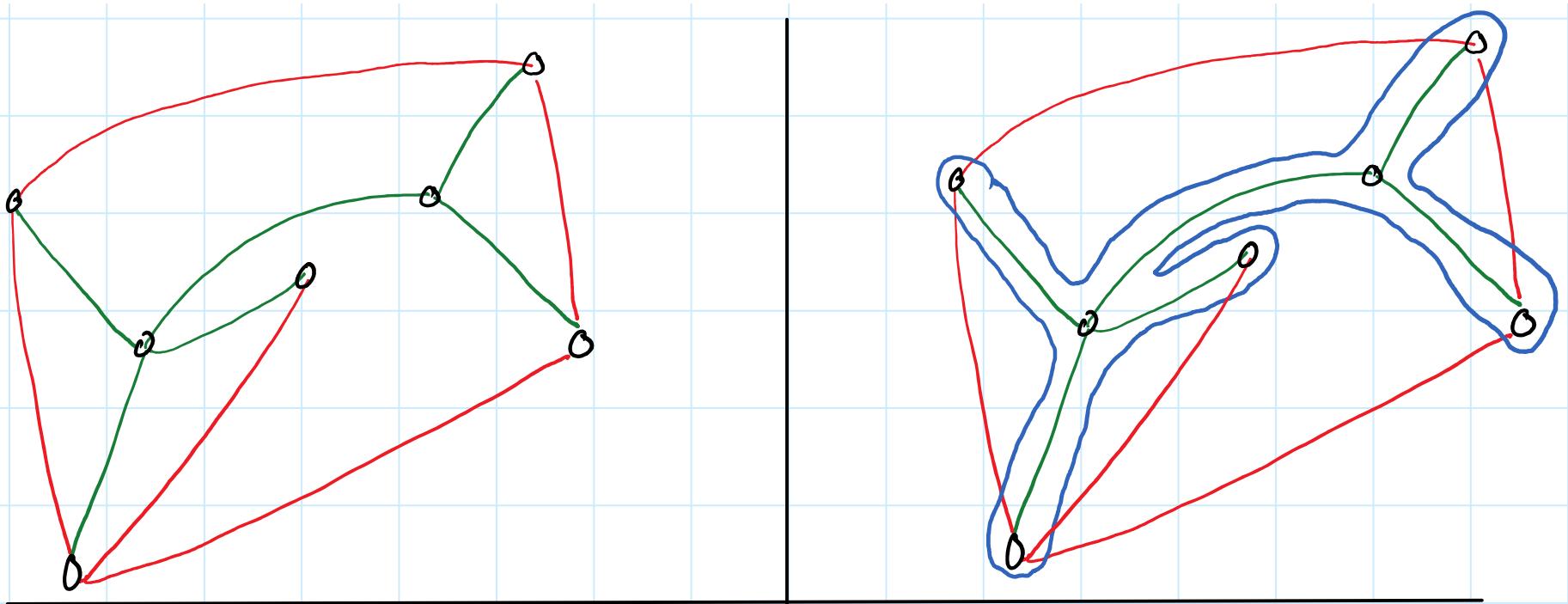
Circle Graph
G

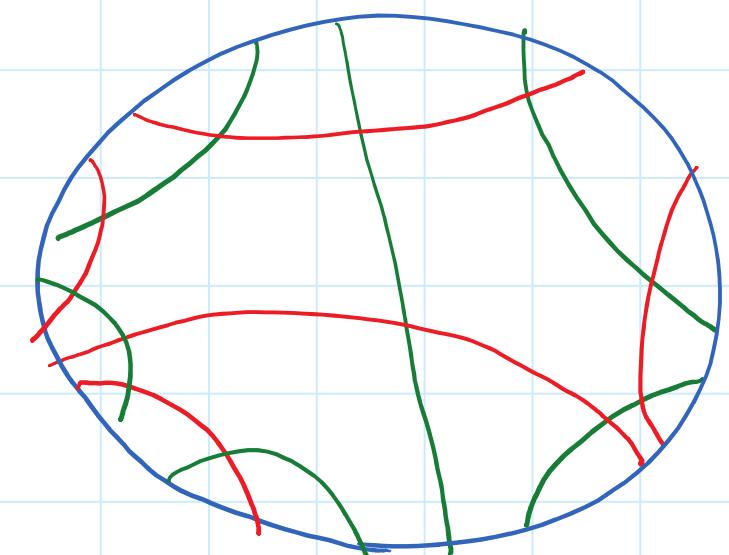
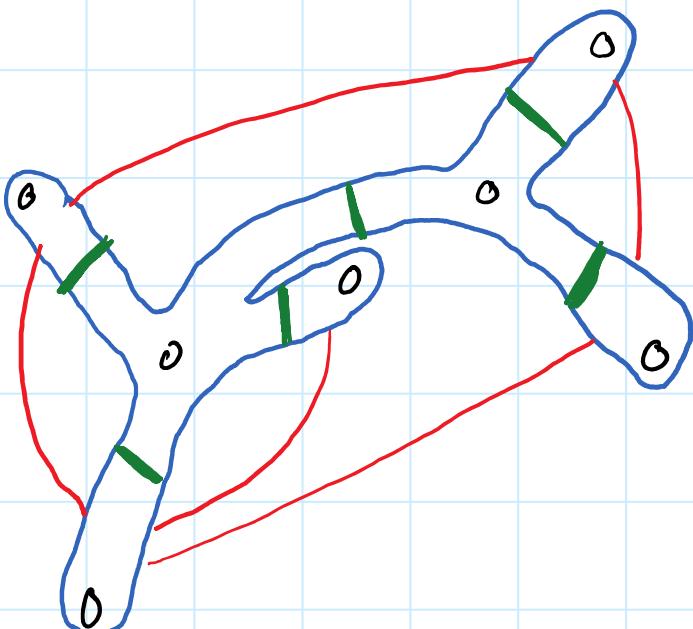
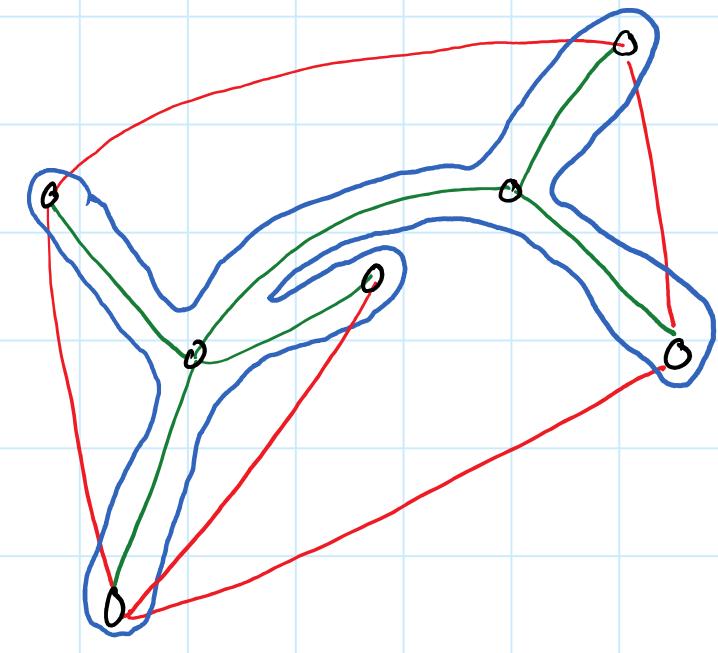
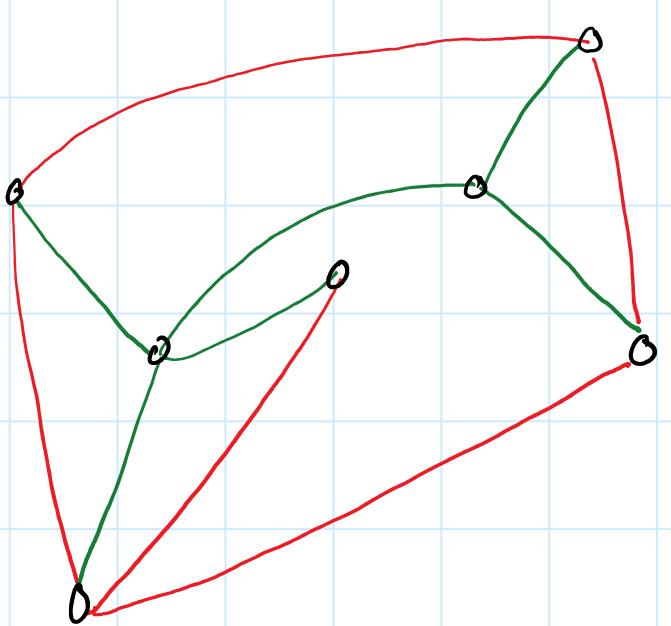


planar graph
green - spanning tree
red - everything else

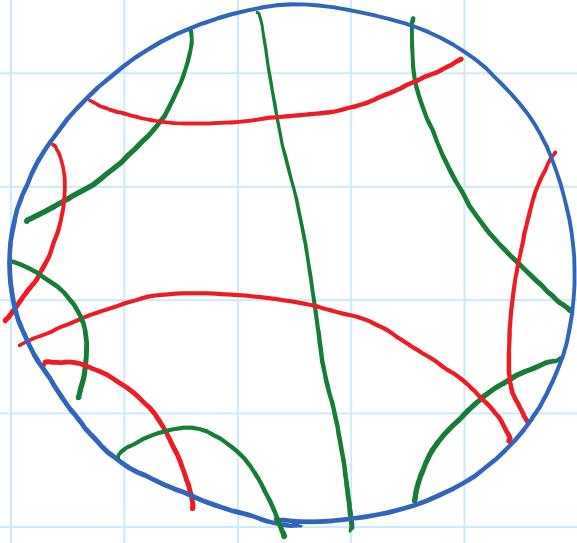






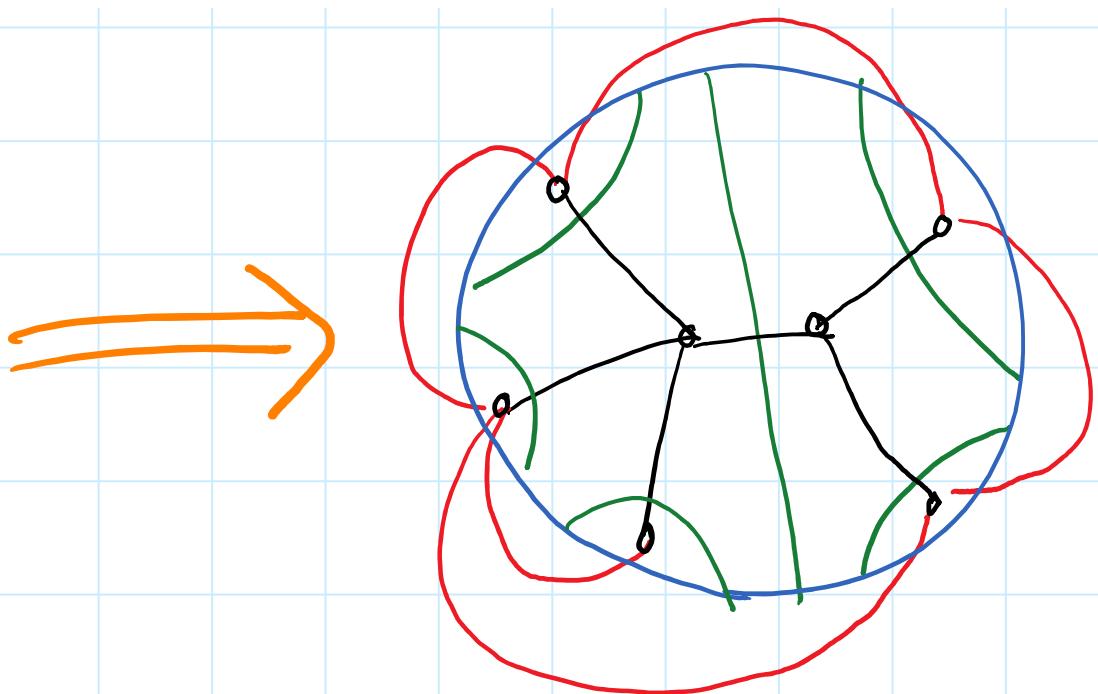
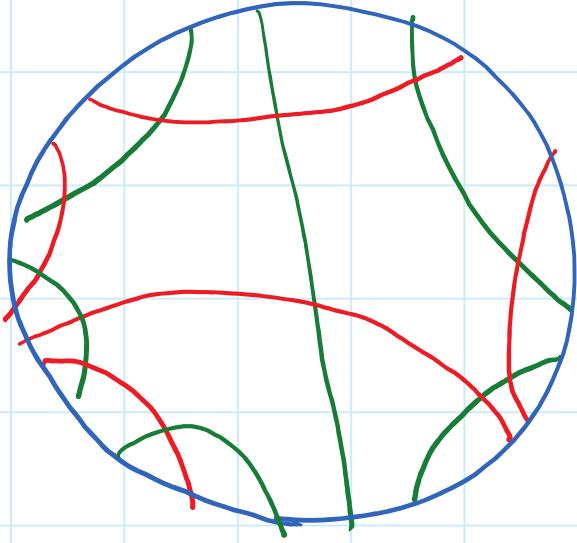


bipartite chord diagram

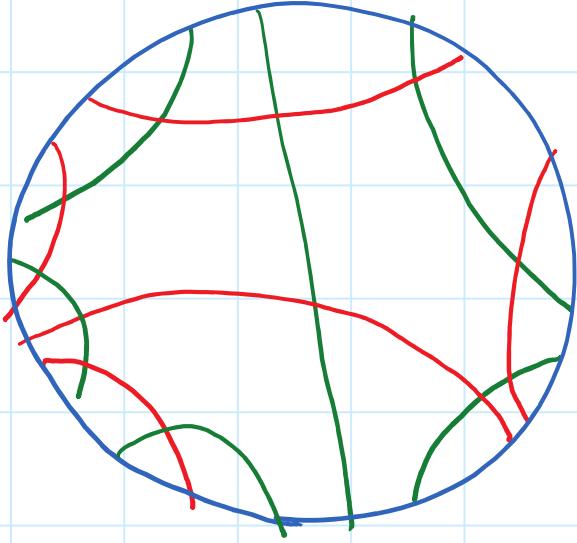


bipartite chord diagram

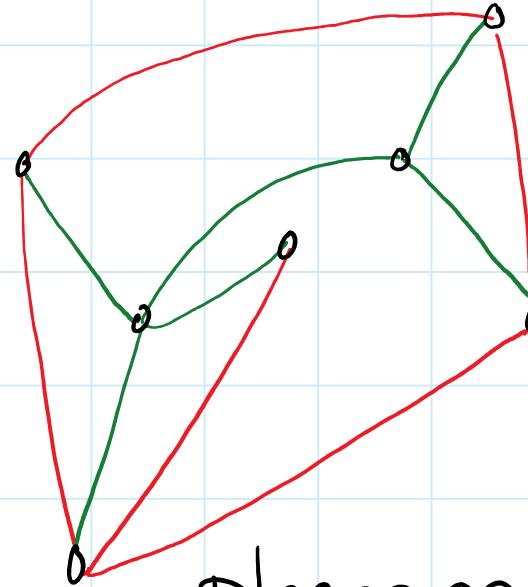
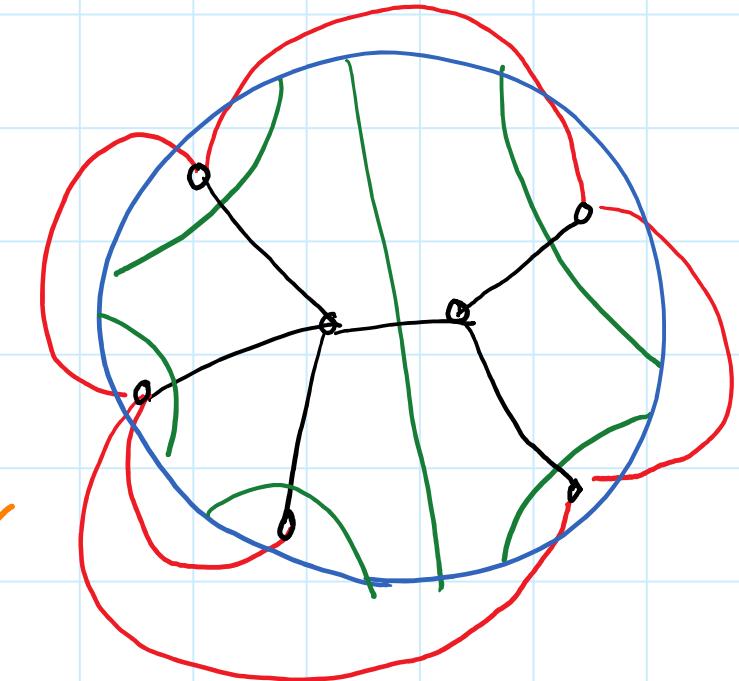
This process is also reversible.



bipartite chord diagram



bipartite chord diagram



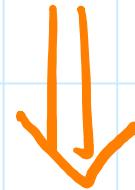
Planar graph

This construction is known as

De Frayessix's Theorem

The important takeaway is that:

characterizations of circle graphs



characterizations of planar graphs

Naji System for G :

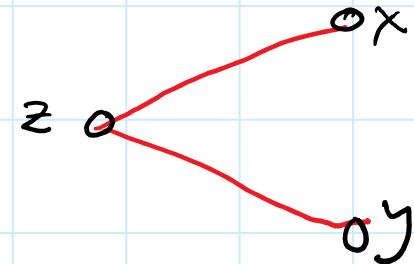
(1) $\beta(x,y) + \beta(y,x) = 1$ for



(2) $\beta(z,x) + \beta(z,y) = 0$ for

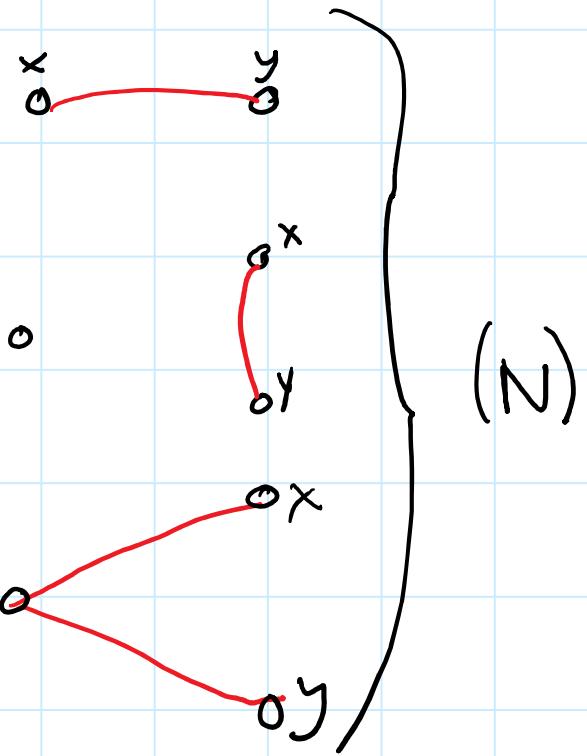


(3) $\beta(z,x) + \beta(z,y)$
 $+ \beta(x,y) + \beta(y,x) = 1$ for



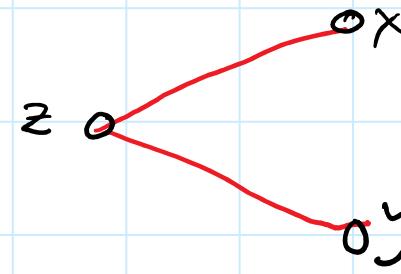
Naji System for G :

(1) $\beta(x,y) + \beta(y,x) = 1$ for



(2) $\beta(z,x) + \beta(x,y) = 0$ for

(3) $\beta(z,x) + \beta(z,y)$
 $+ \beta(x,y) + \beta(y,z) = 1$ for



Theorem (Naji) G is a circle graph $\iff (N)$
has a solution over $GF(2)$.

An interlude into complexity:

- (N) has n^2 variables
 $\sim n^3$ equations

$$\begin{bmatrix} n^3 \\ \vdots \\ \vdots \\ \vdots \\ n^3 \end{bmatrix} = \begin{bmatrix} n^2 \\ \vdots \\ \vdots \\ \vdots \\ n^2 \end{bmatrix}$$

The diagram illustrates a system of equations where the left side is a column vector of length n^3 , and the right side is a column vector of length n^2 . The left side is partitioned into n^3 rows, each containing n^2 entries, represented by red 'x' marks. The right side is partitioned into n^2 rows, each containing n^2 entries, also represented by red 'x' marks. A red bracket on the left side groups the first n^2 entries of each row, and another bracket on the right side groups the first n^2 entries of each row, indicating a correspondence between the two sides of the equation.

An interlude into complexity:

- (N) has n^2 variables
 $\sim n^3$ equations

$$\begin{bmatrix} n^3 \\ \vdots \\ n^3 \end{bmatrix} = \begin{bmatrix} n^2 \\ \vdots \\ n^2 \end{bmatrix}$$

A diagram showing two vectors side-by-side. The left vector is labeled n^3 above its first entry and has red vertical bars on its sides. The right vector is labeled n^2 above its first entry and has red vertical bars on its sides. An equals sign is placed between the two vectors.

- Gaussian (column) elimination can be done in $O(n^5)$ time for each column.

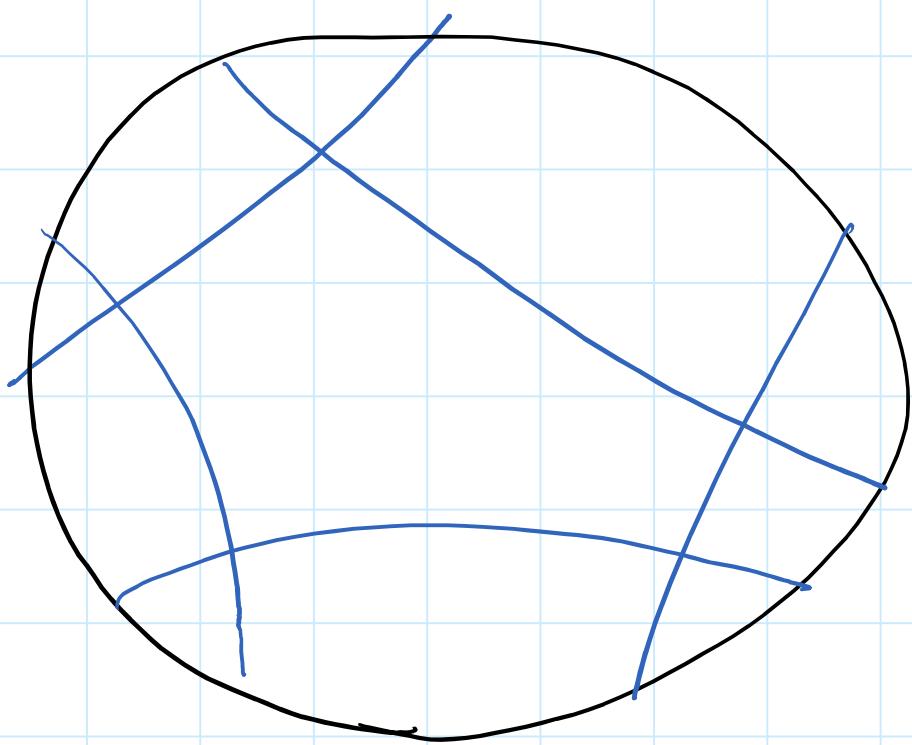
An interlude into complexity:

- (N) has n^2 variables
 $\sim n^3$ equations

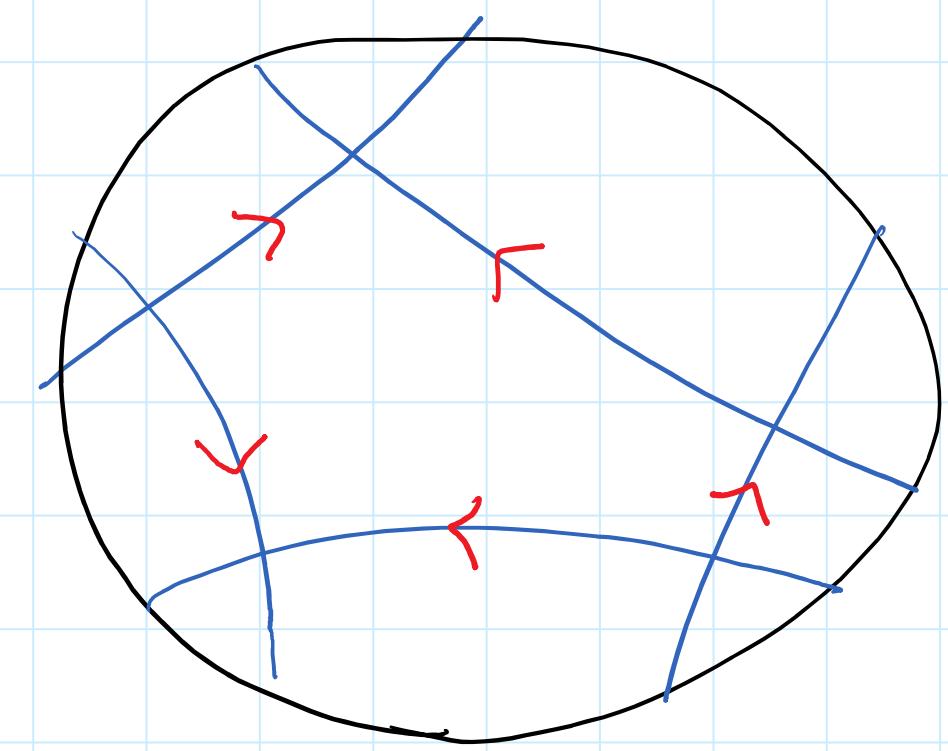
$$\begin{bmatrix} \vdots & & & \\ \vdots & \ddots & & \\ \vdots & & \ddots & \\ \vdots & & & \ddots \end{bmatrix}^{n^3} = \begin{bmatrix} \vdots & & & \\ \vdots & \ddots & & \\ \vdots & & \ddots & \\ \vdots & & & \ddots \end{bmatrix}^{n^2}$$

- Gaussian (column) elimination can be done in $O(n^5)$ time for each column.
- So $O(n^7)$ algorithm for recognizing circle graphs.

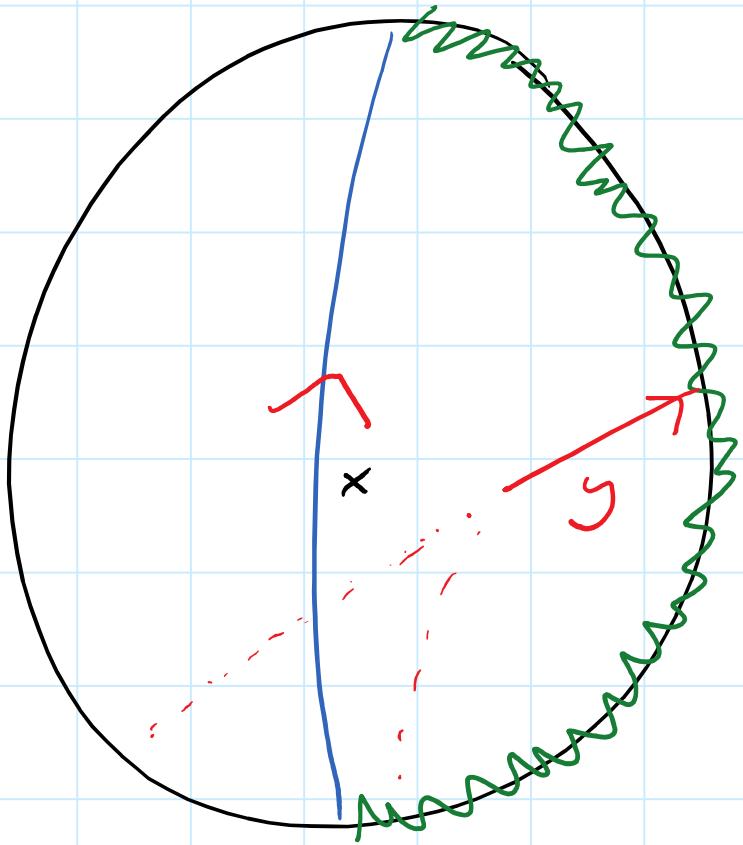
Circle Graph \Rightarrow Solution



chord diagram



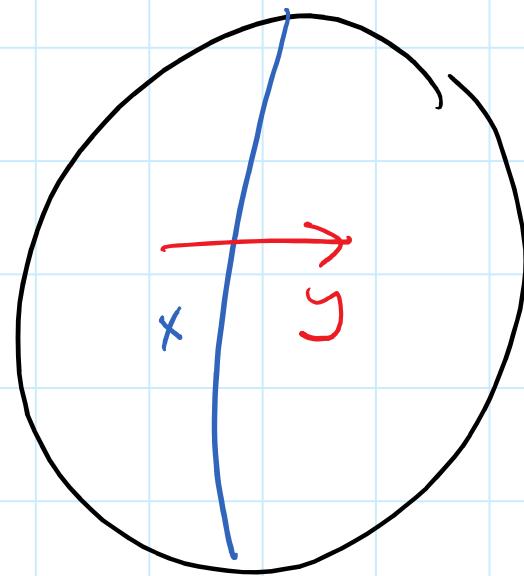
+orientations



Take $\beta(x,y) = 0$ if the head of y is to the right of x .

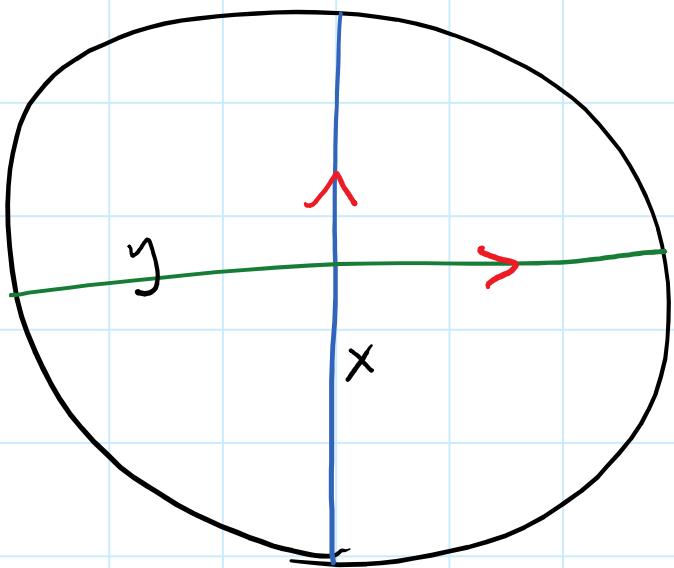
Notation: $\beta(\ell)$ for such solutions.

Aside: will often draw
as:



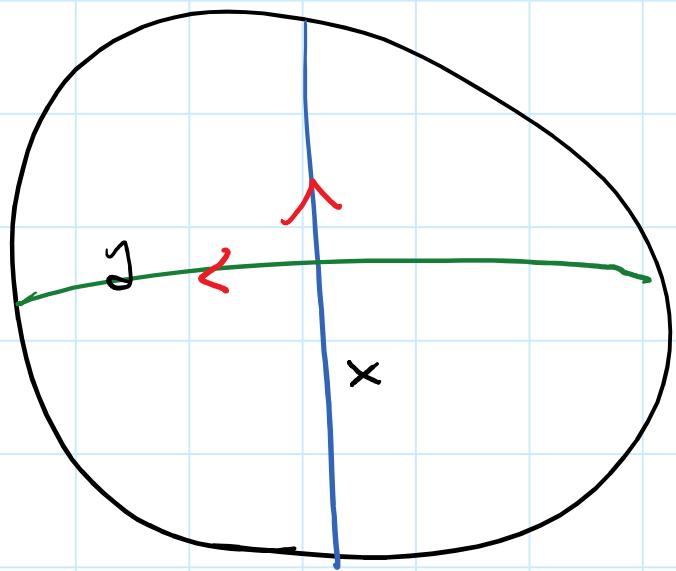
Equation (1): $x \circ \circ y$

$$\beta(x,y) + \beta(y,x) = 1$$



$$\beta(x,y) = 1$$

$$\beta(y,x) = 0$$



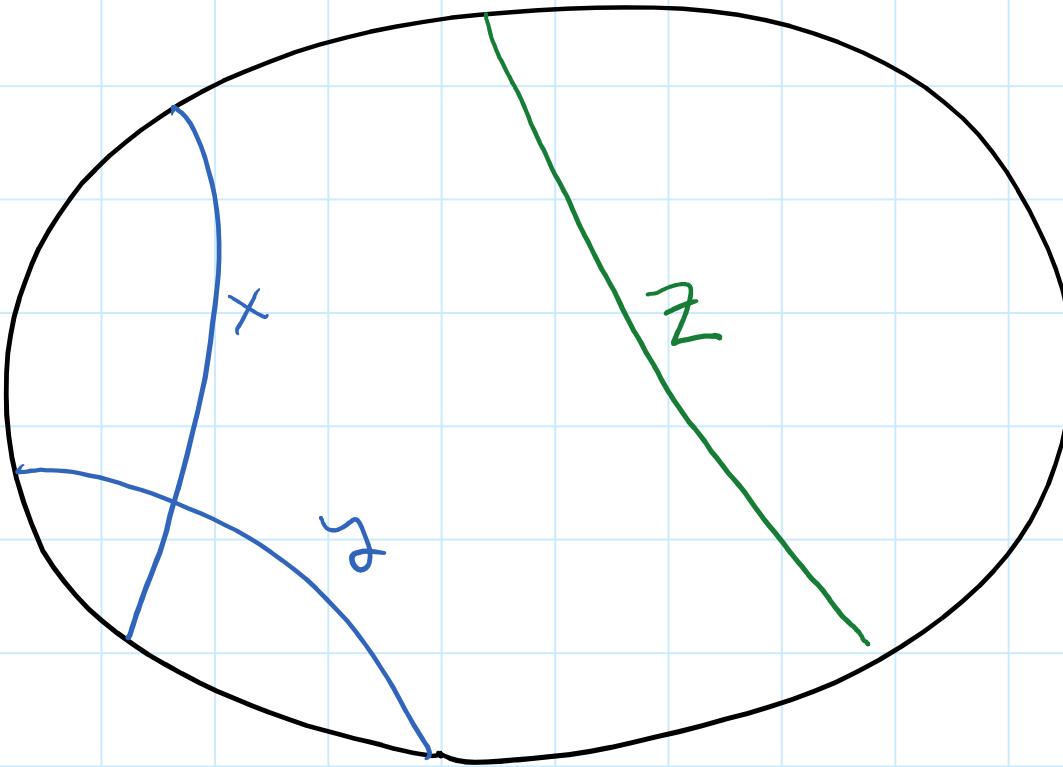
$$\beta(y,x) = 1$$

$$\beta(x,y) = 0$$

Equation (2): z^0

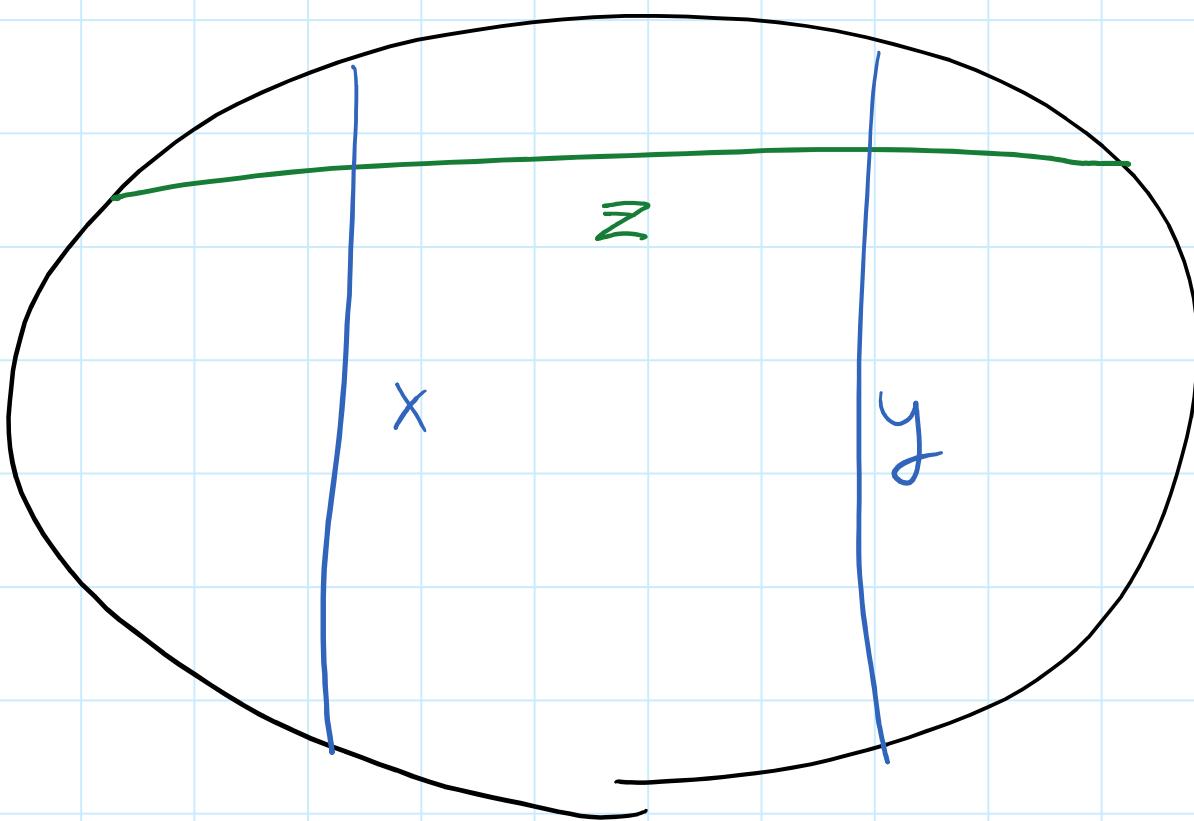
$$\beta(z, x) + \beta(z, y) = 0$$

cx
 cy



Equation (3): z^o

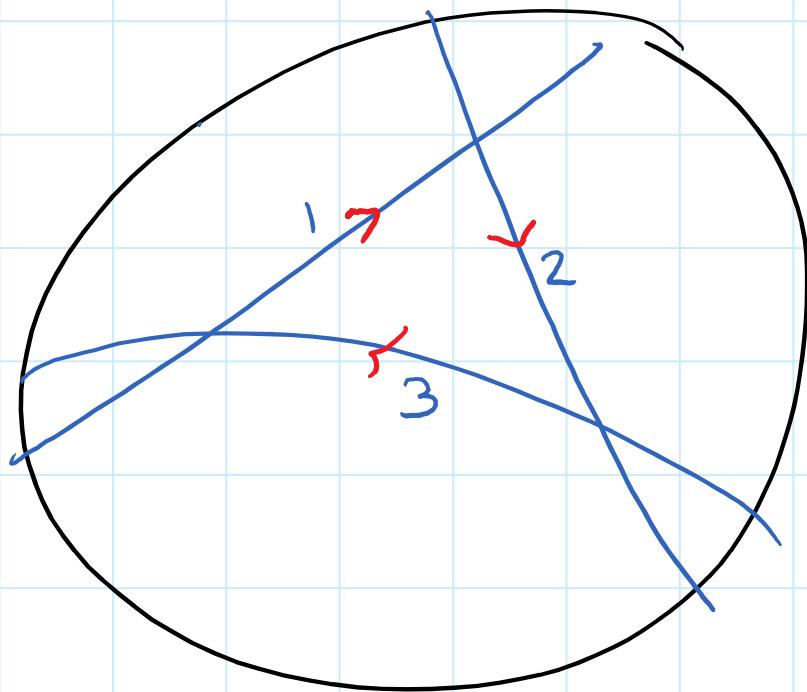
$$\beta(z, x) + \beta(z, y) + \beta(x, y) + \beta(y, z) = 1$$

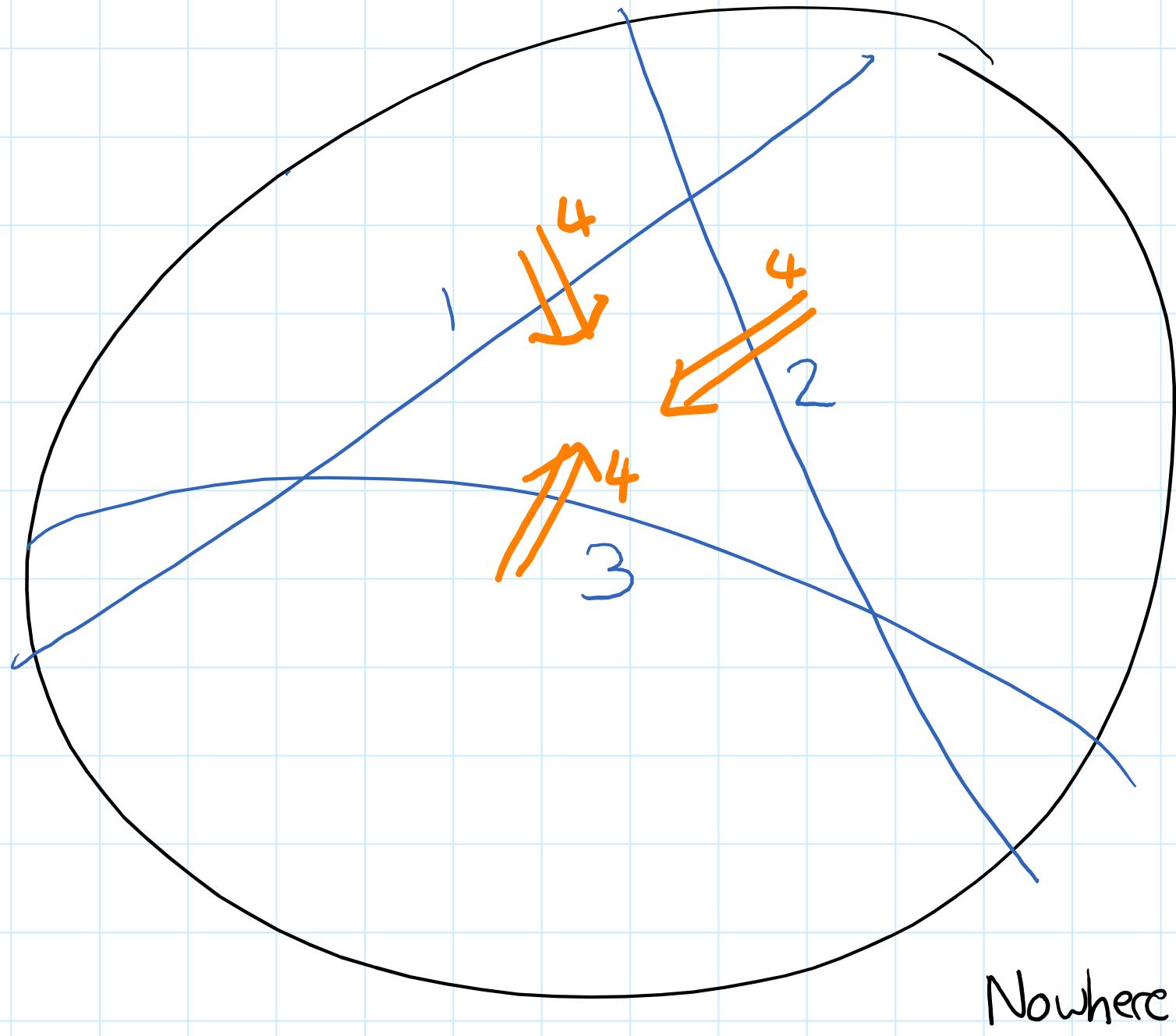


Not all Naji solutions are so natural.

Consider the following solution to K_4 :

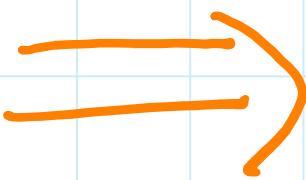
$$\begin{array}{l|l} \beta(1,2) = \beta(2,3) = \beta(3,1) = 0 & \text{everything else } 1. \\ \beta(1,4) = \beta(2,4) = \beta(3,4) = 0 \end{array}$$





Nowhere to put 4.

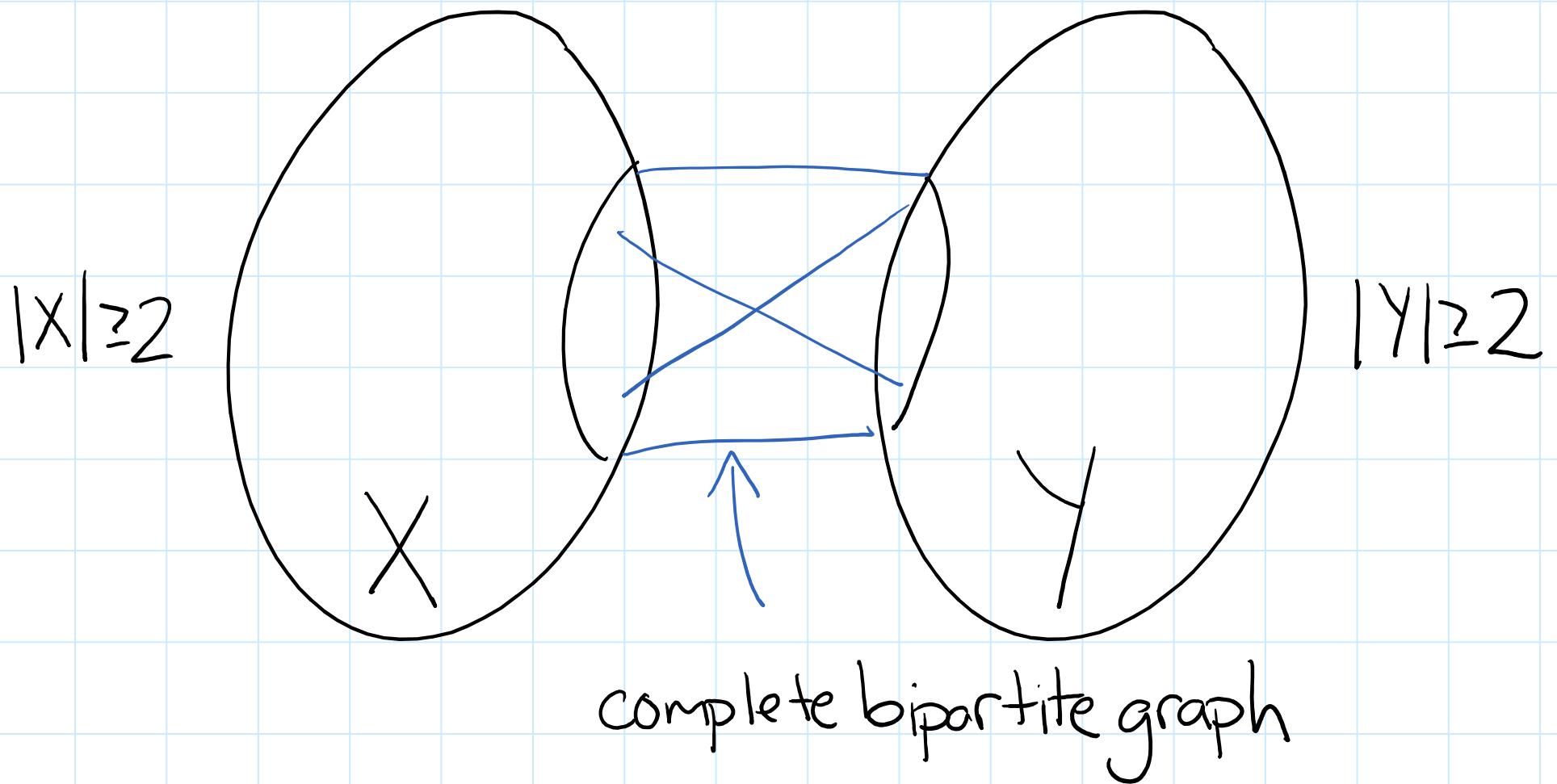
$$A\bar{x} = A\bar{y} = A\bar{z} = b$$

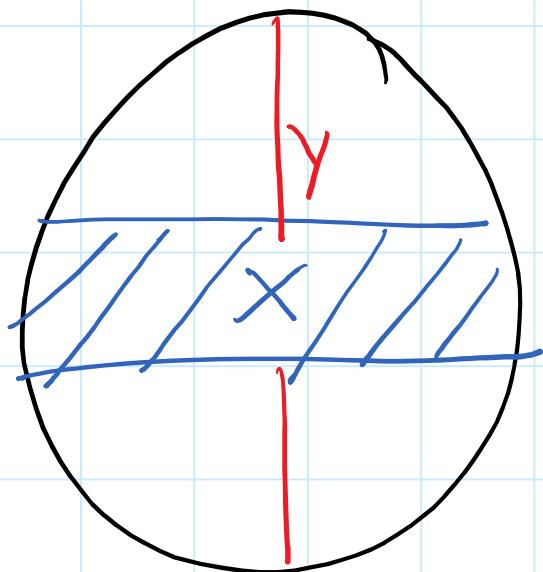
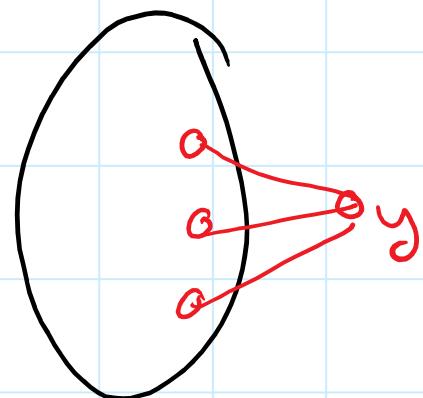


$$A(\bar{x} + \bar{y} + \bar{z}) = b$$

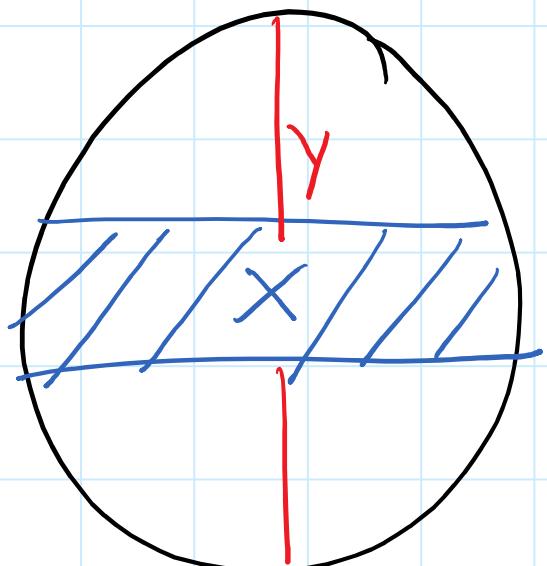
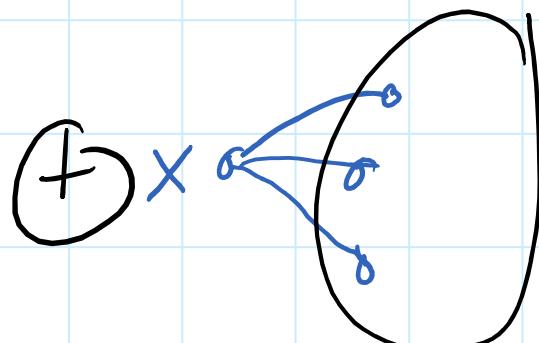
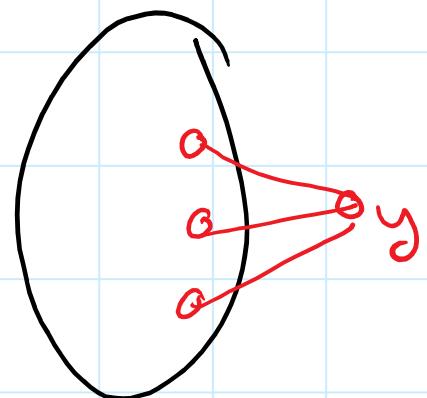
Can combine "unrelated" Naji solutions!

Splits \Rightarrow Inequivalent Diagrams

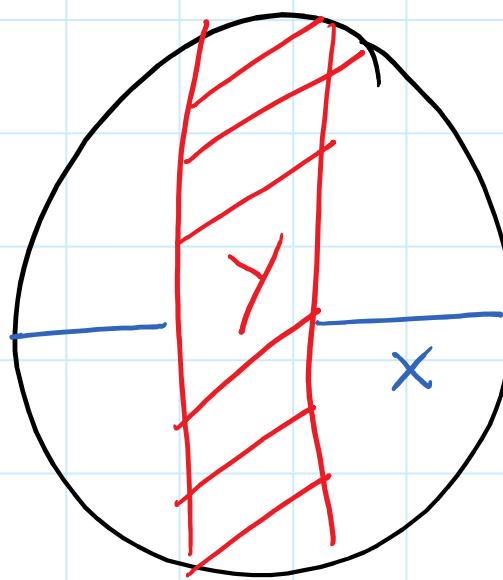




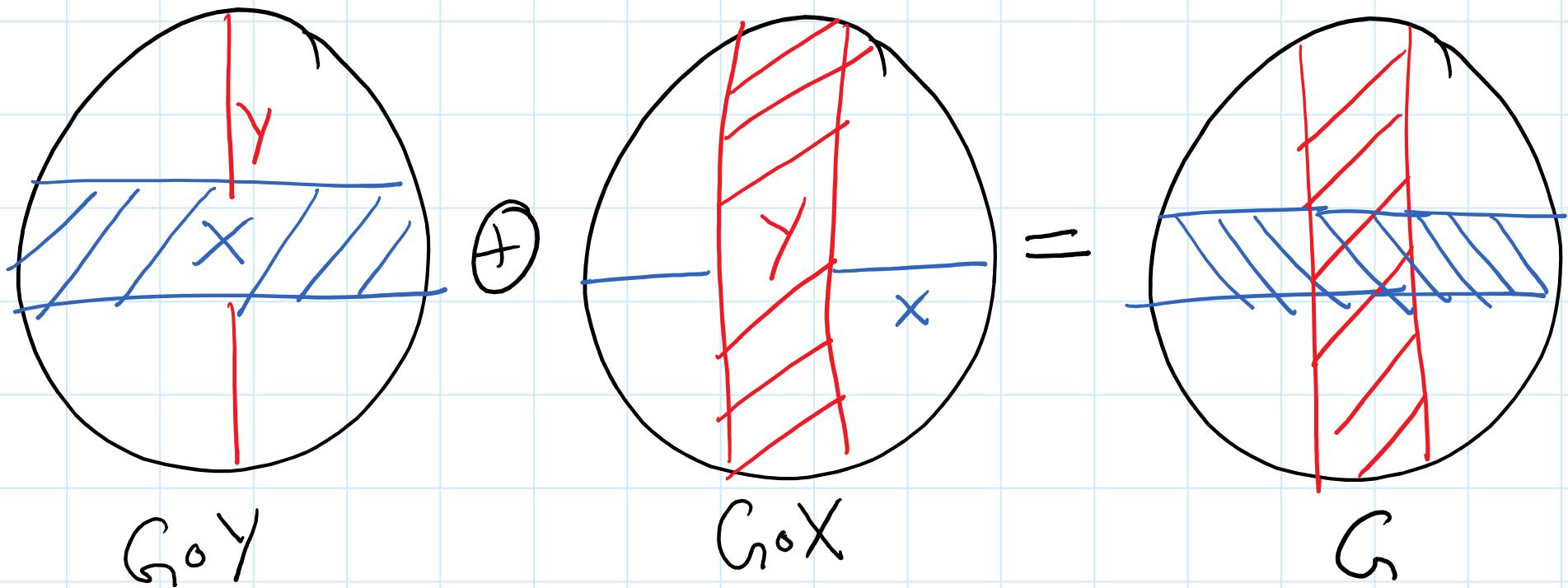
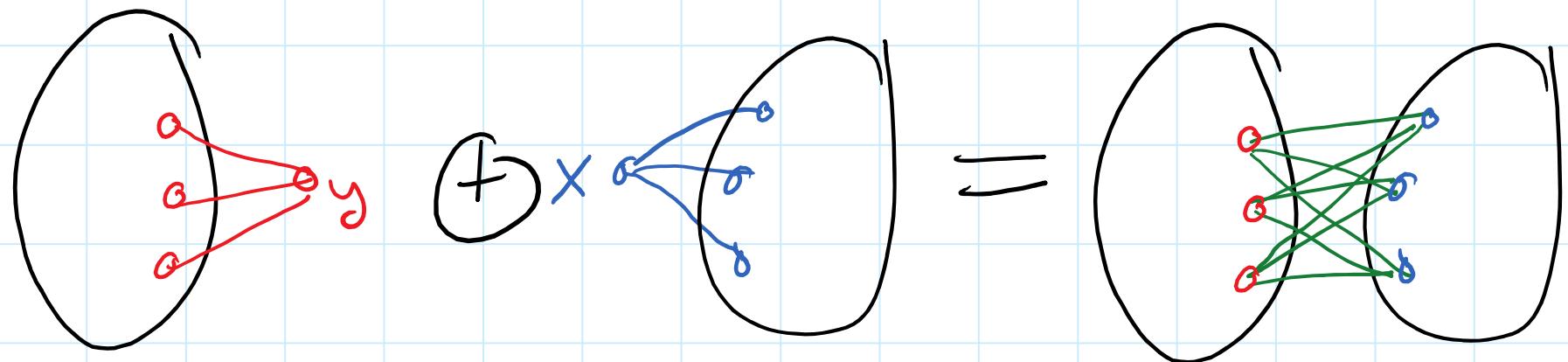
$G \circ Y$



$G_0 Y$



$G_0 X$

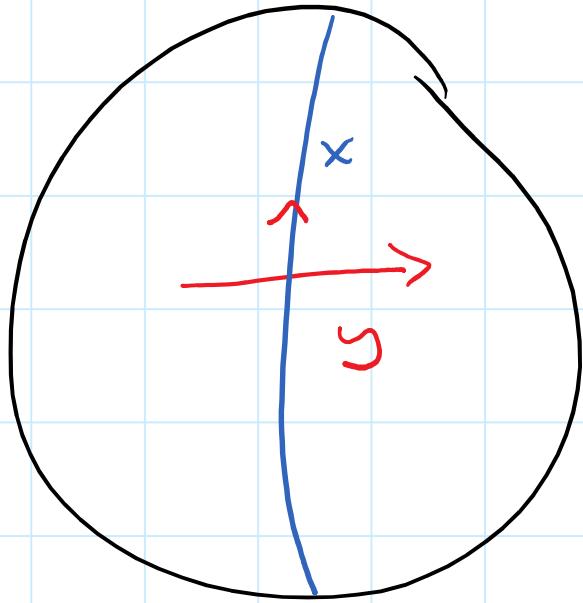


Theorem (G,L)

Let β be a Naji solution for G . Then either:

(1): $\beta = \beta(e)$

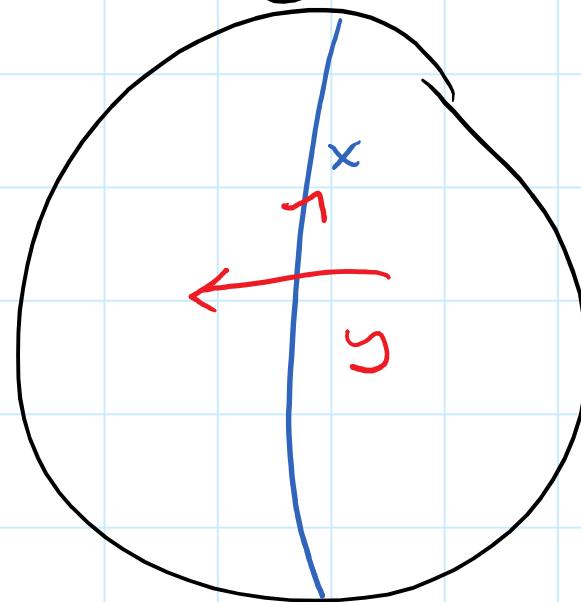
(2): G has a split.



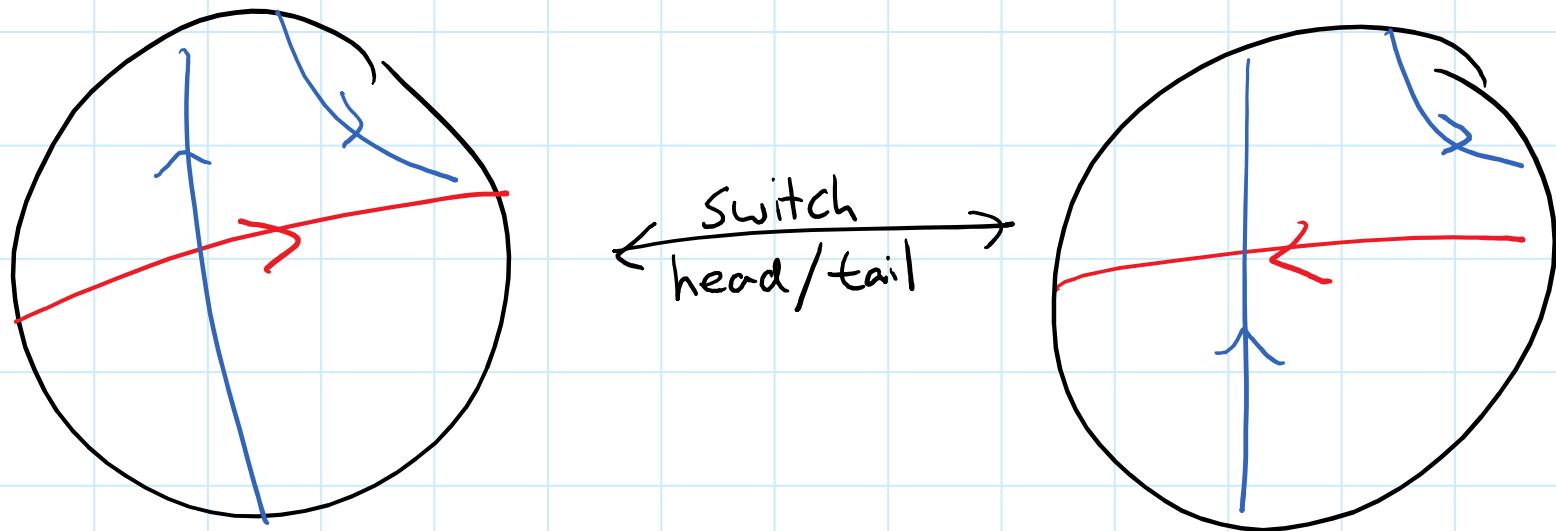
$$B(x,y) = 0$$

Each chord knows
where the head of
the next chord
should go.

$$B(x,y) = 1$$



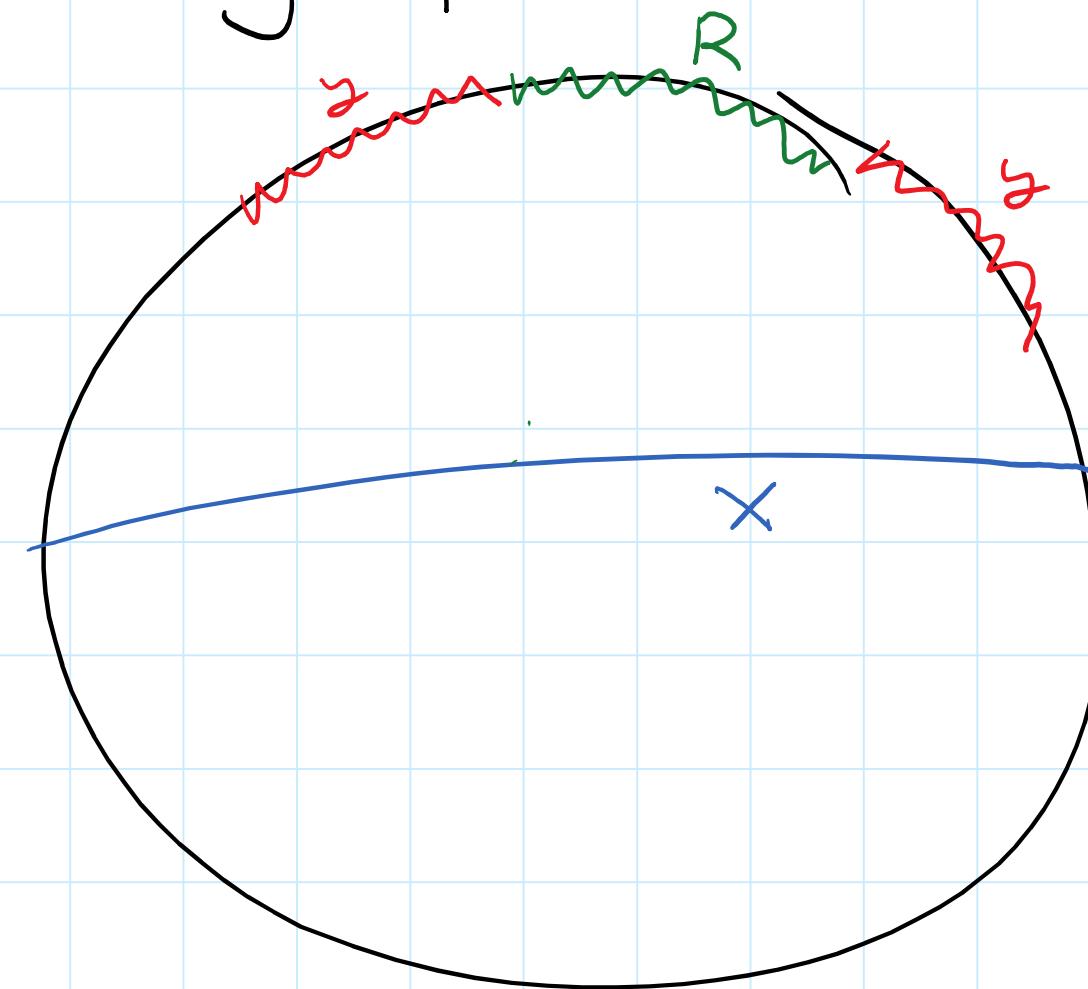
Moreover we can easily modify solution
to find out where the tail should go. (Gasse)



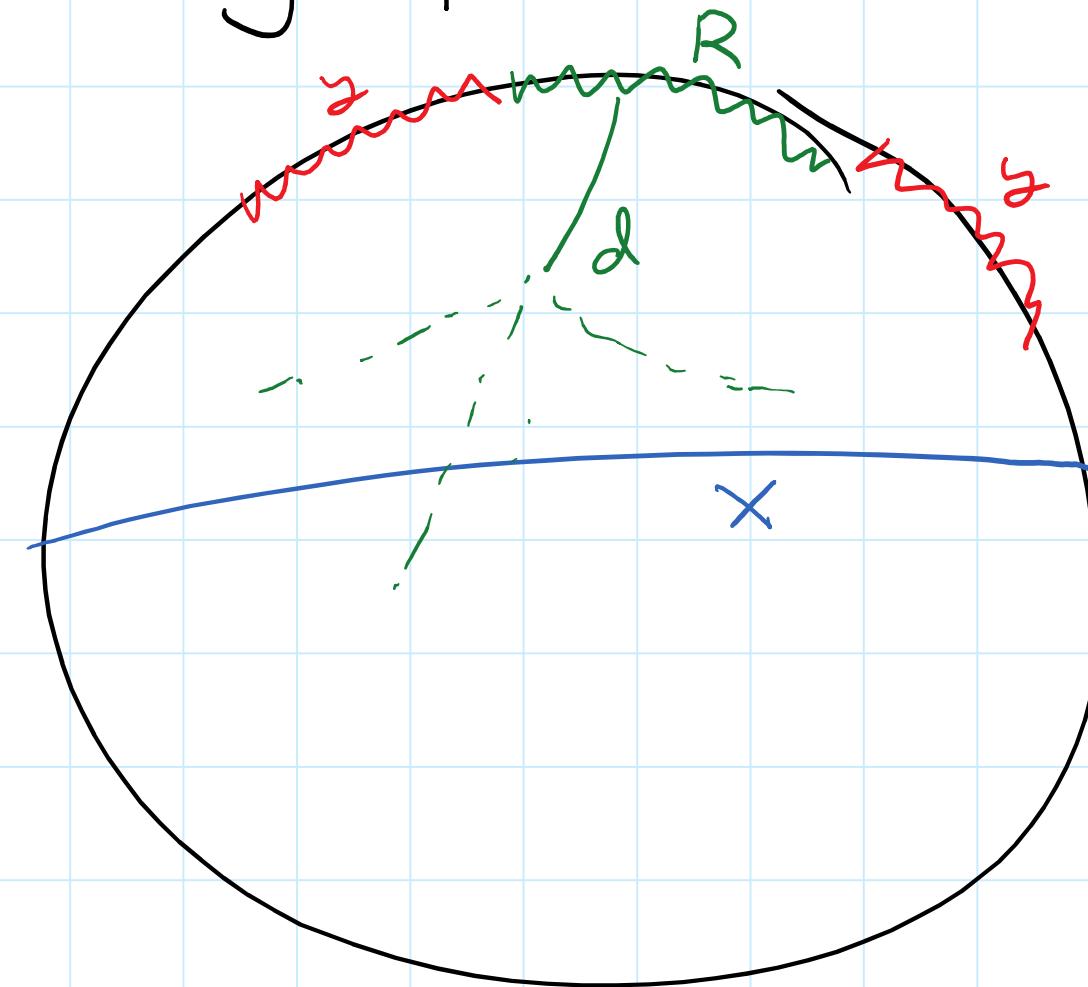
$$\beta(e) = \gamma(v) + \beta(e')$$

where $\gamma_v(x,y)=1$ if and only if $x=v$
or $y=v$ and $\{x,y\} \in E$.

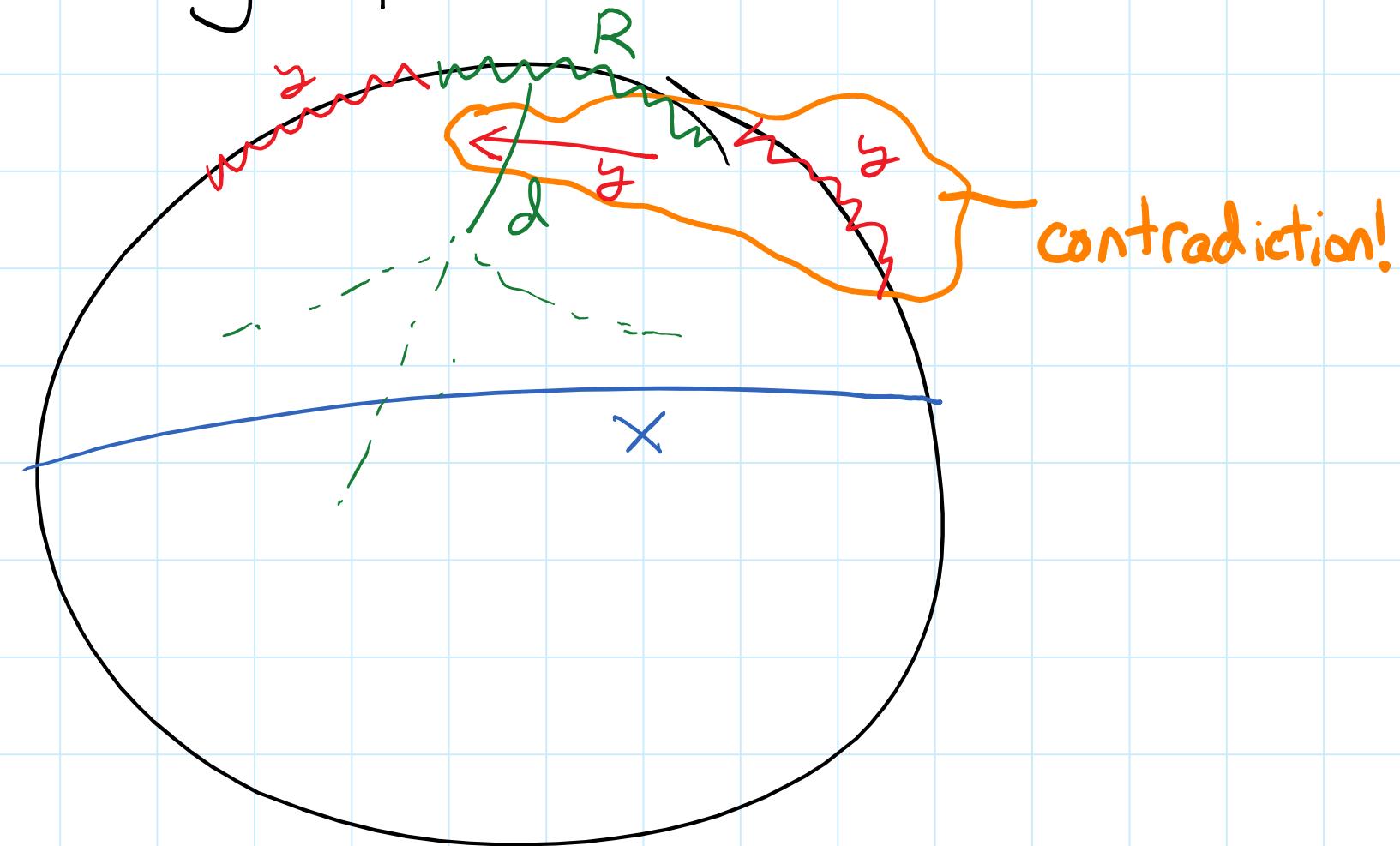
So long as we place the chords in a connected manner there is at most one way to place a new chord.



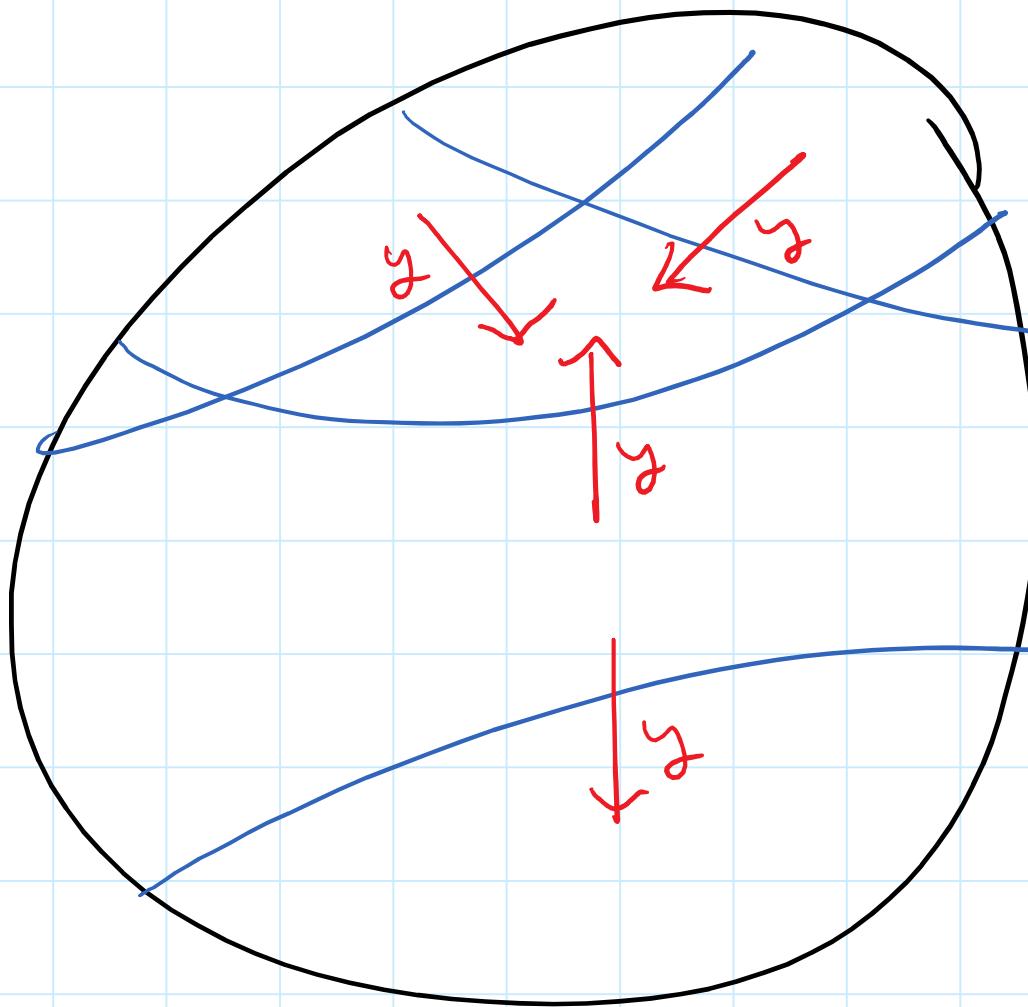
So long as we place the chords
in a connected manner there is at
most one way to place a new chord



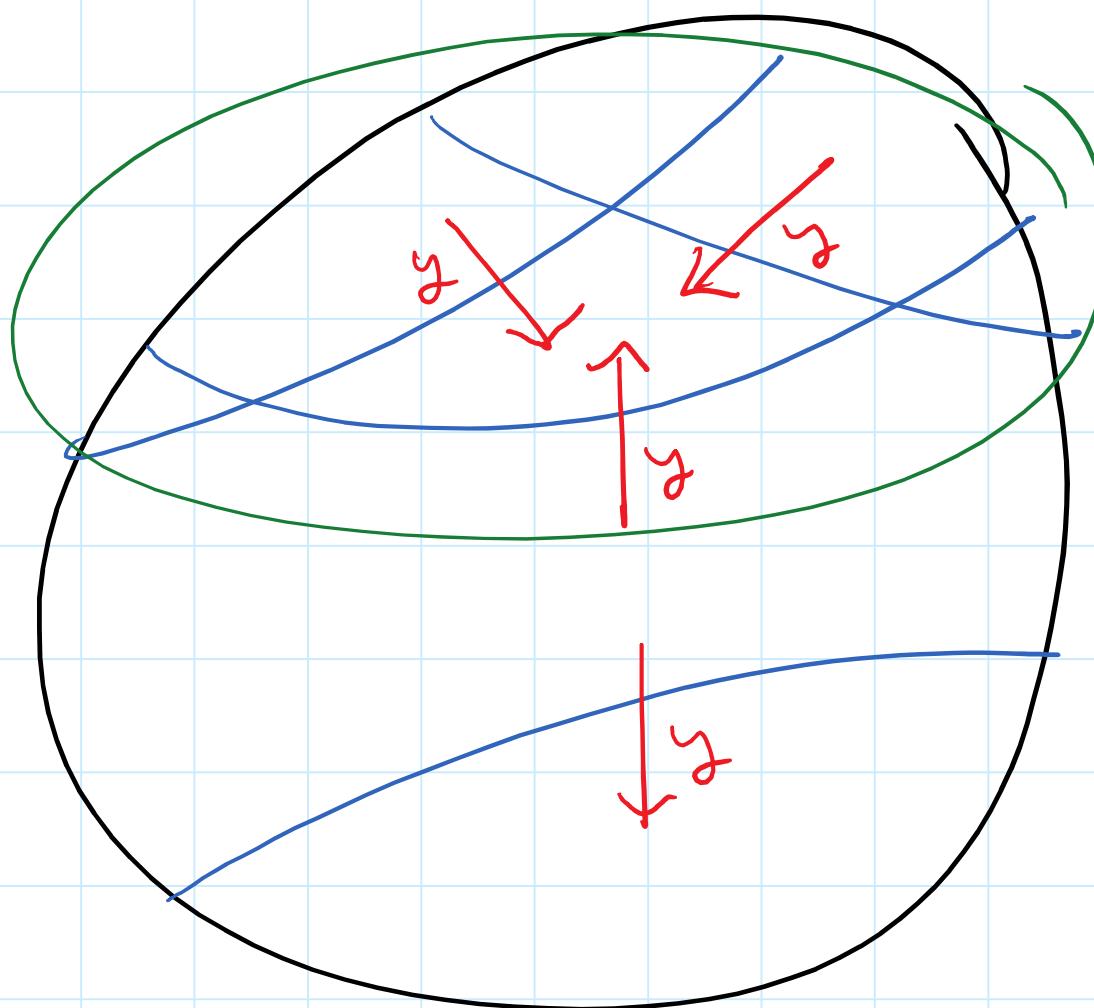
So long as we place the chords in a connected manner there is at most one way to place a new chord.



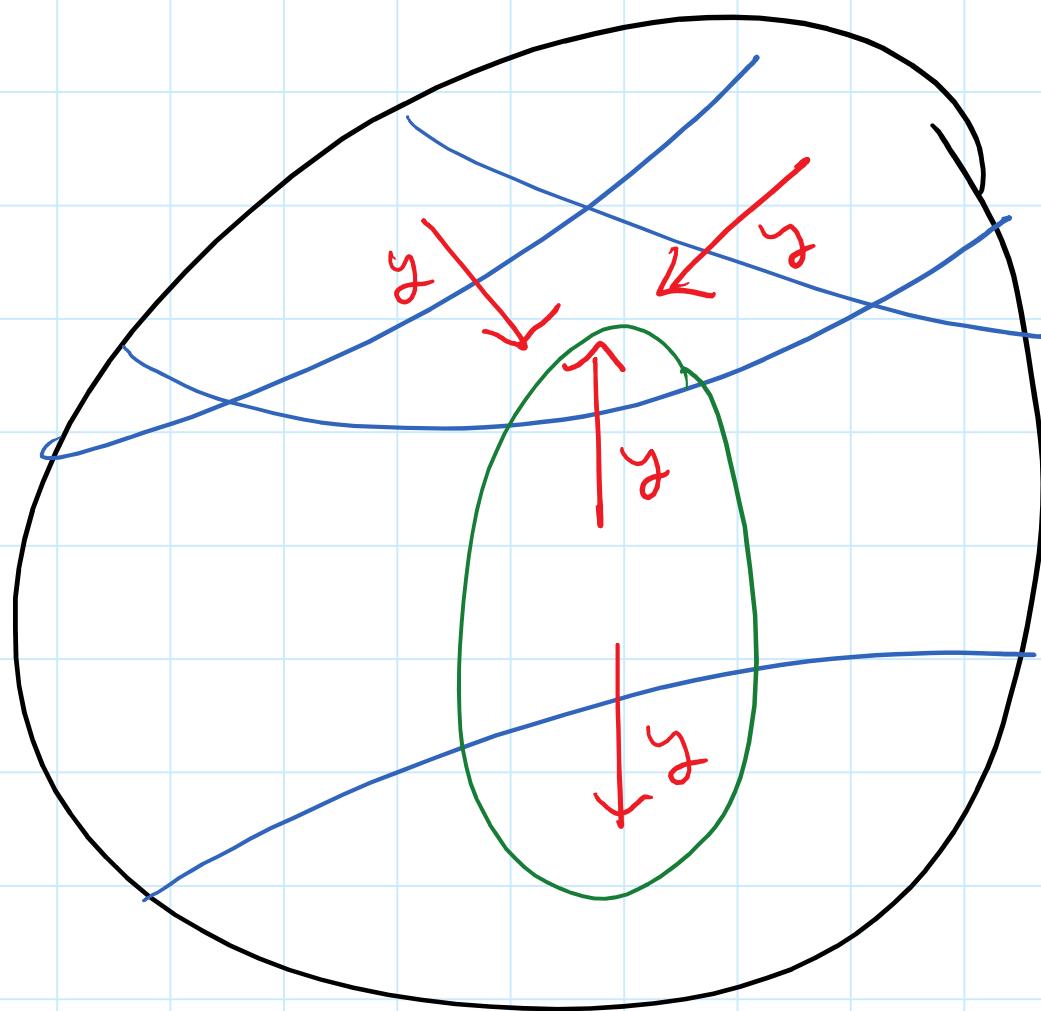
If something goes wrong:



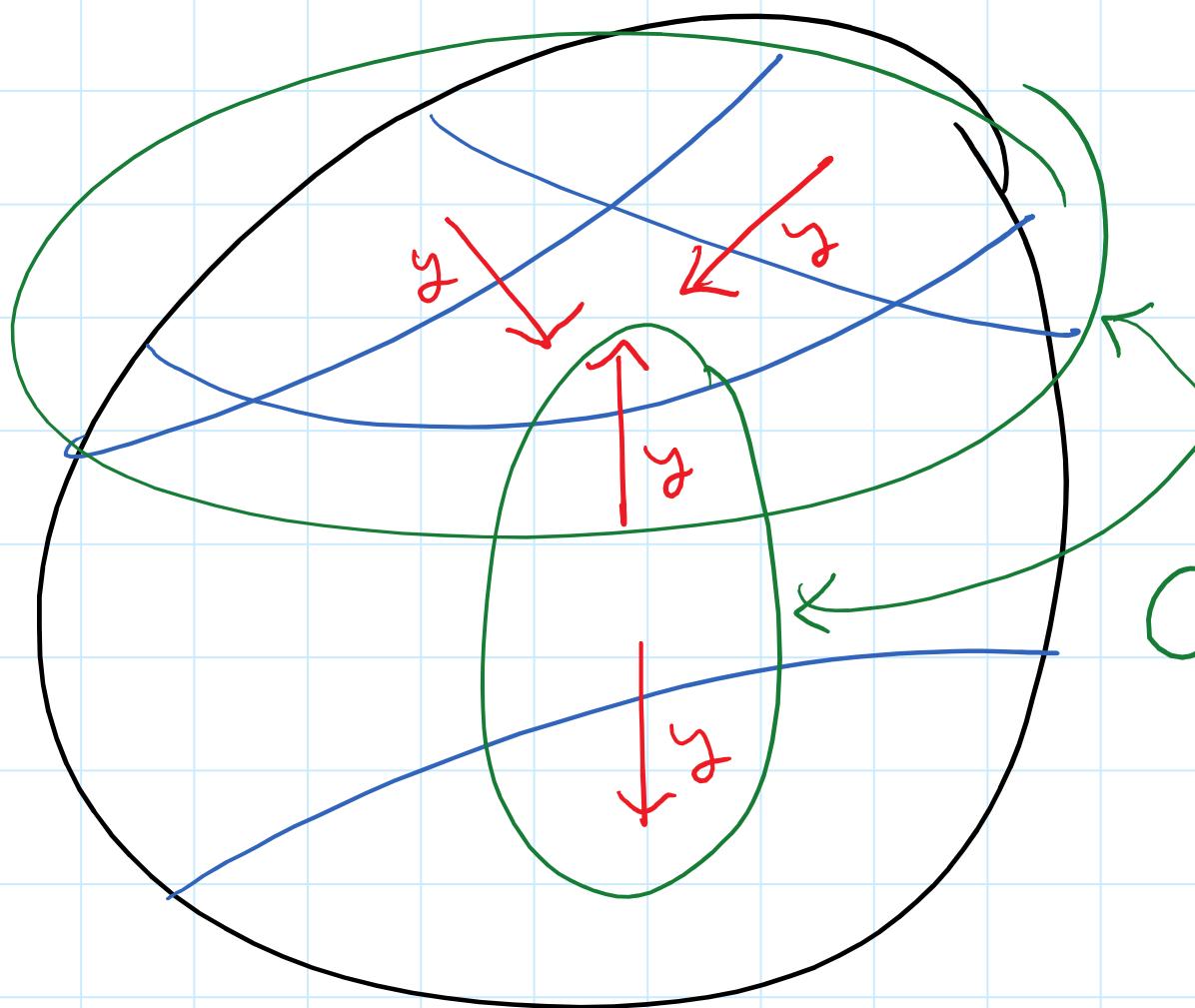
If something goes wrong:



If something goes wrong:

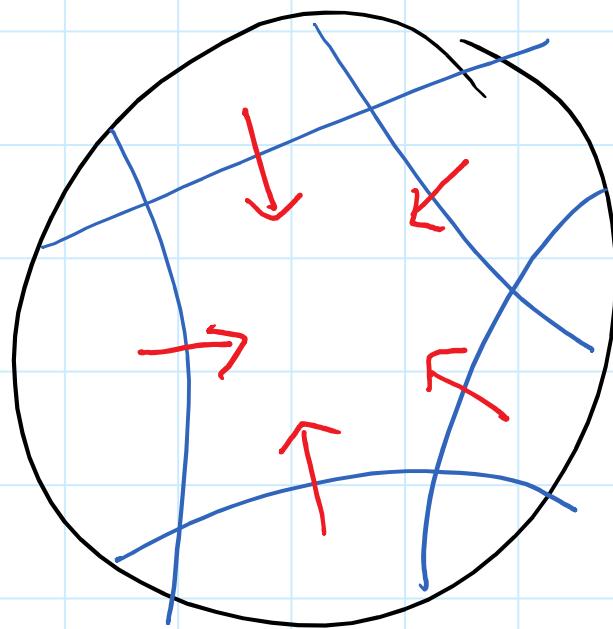


If something goes wrong:

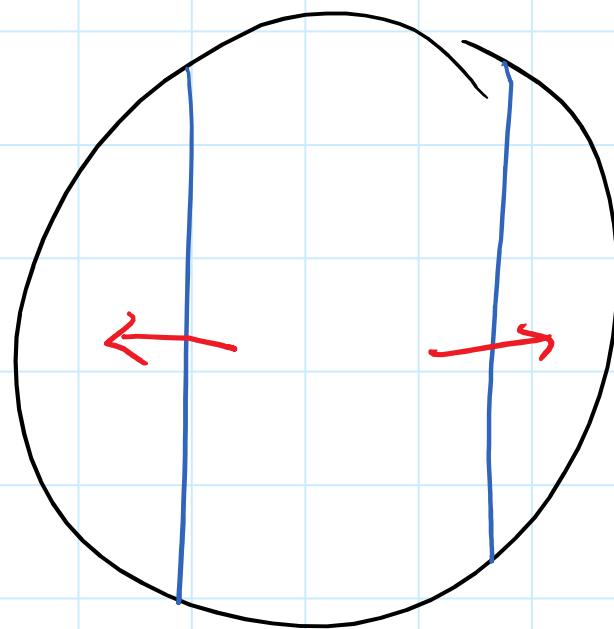


Chords
with
Contradictions

Now if there are chords with contradictions
we can find a minimal contradictory set:

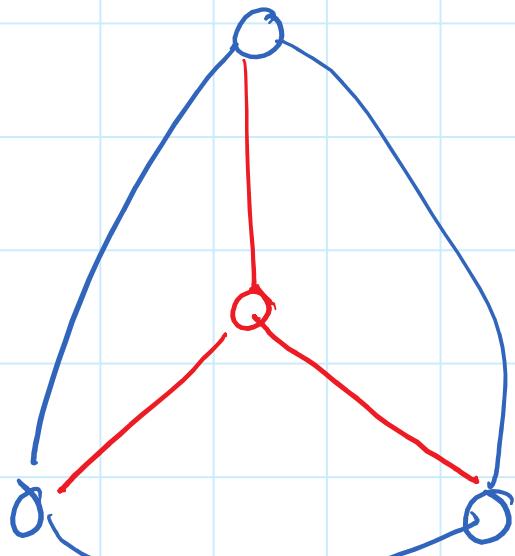
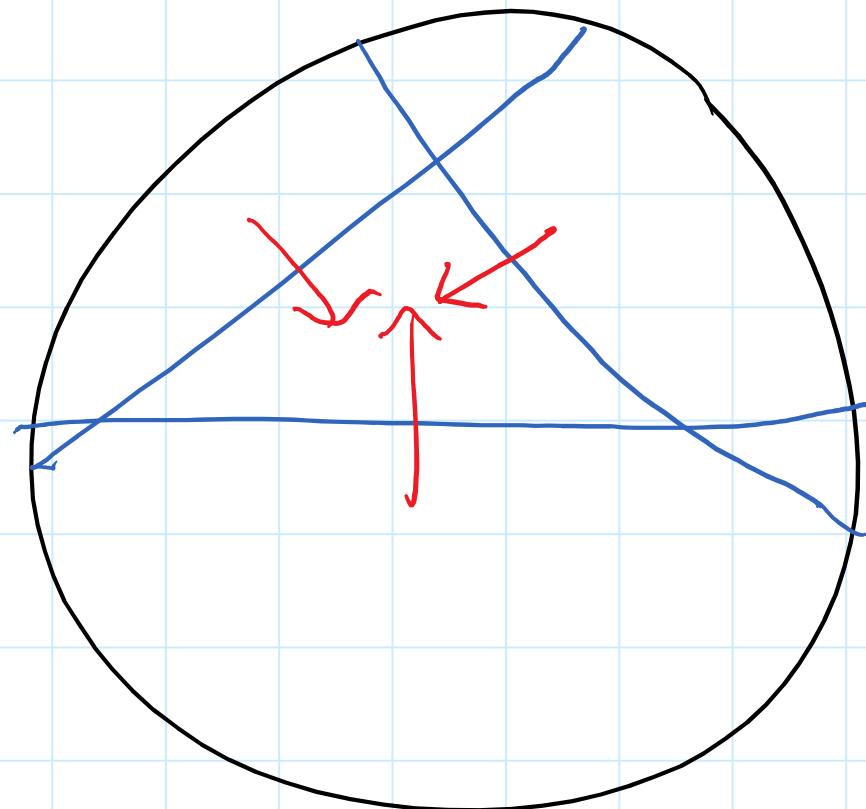


Induced C_k ,
 $k \geq 3$



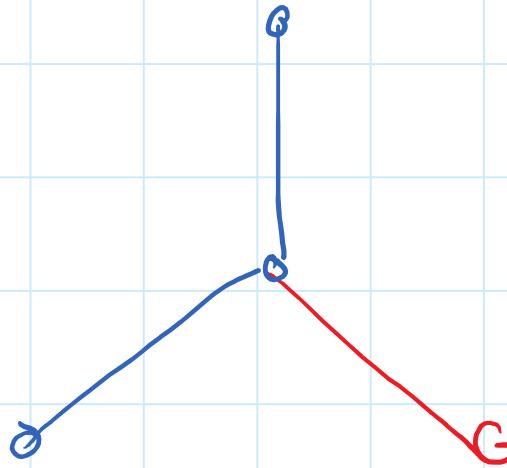
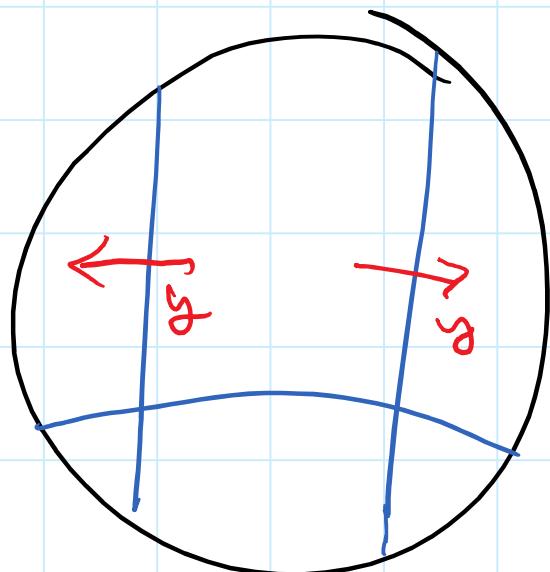
Two non-adjacent
vertices

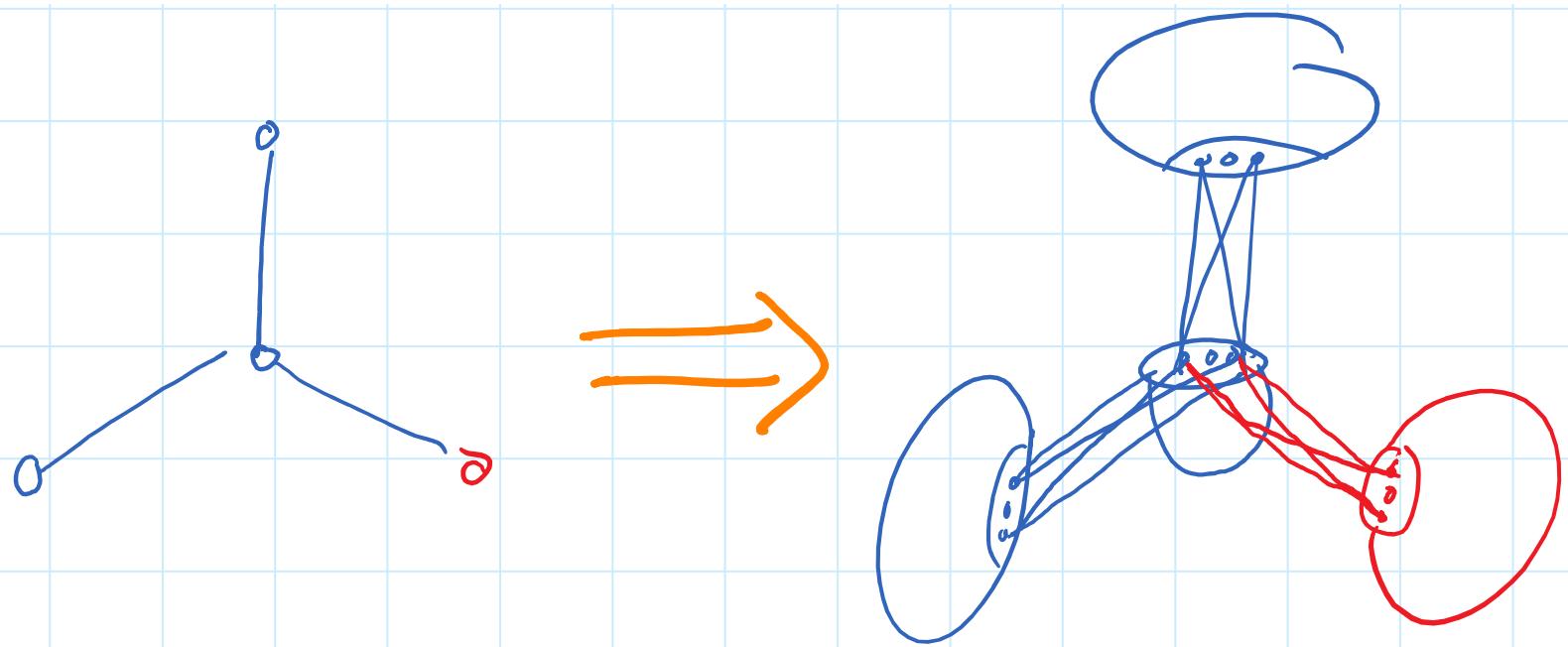
- If $k \geq 5$ then G violates (N)
- So $C_k \hookrightarrow C_3 + \text{axle} = K_4$

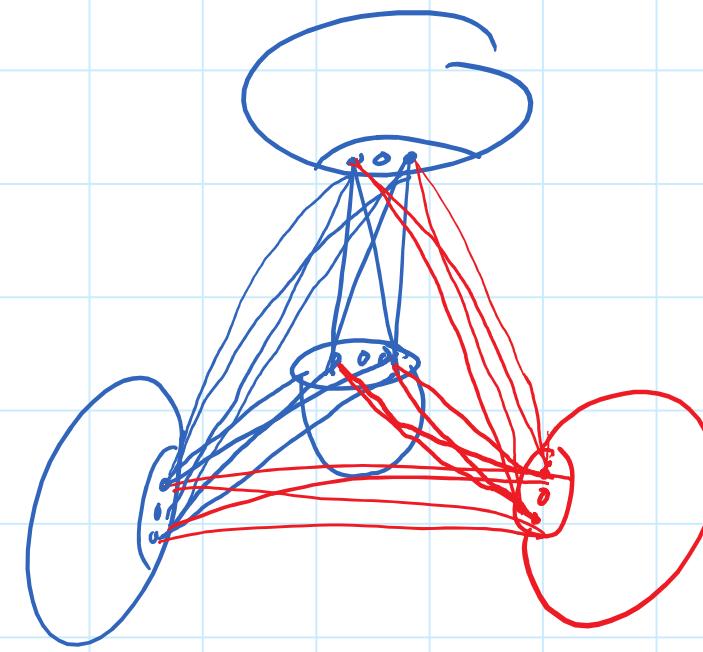
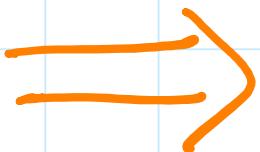
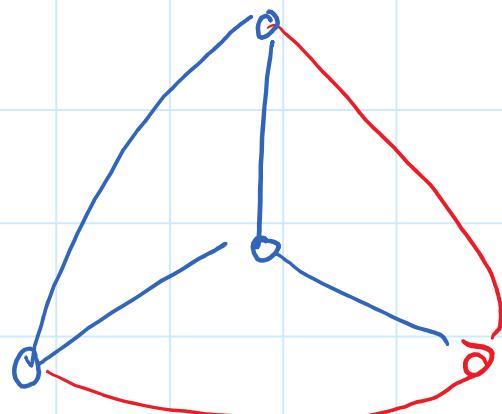
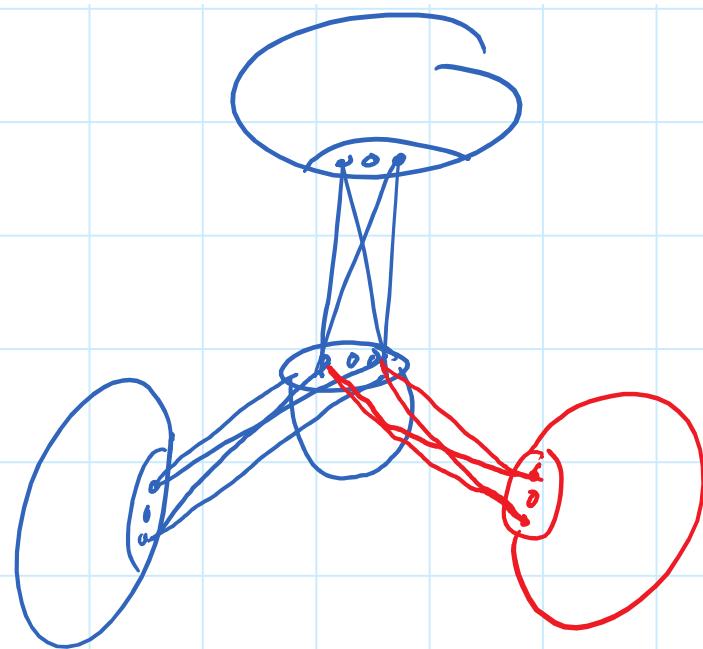
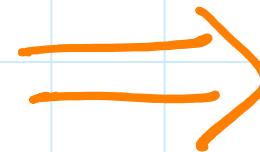
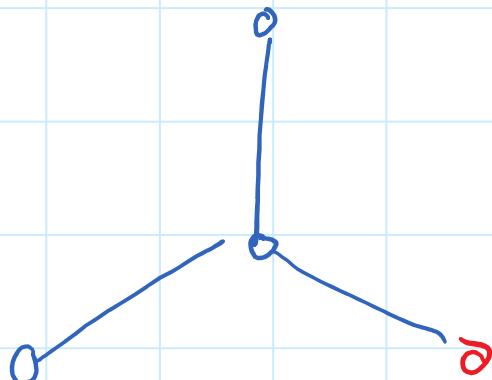


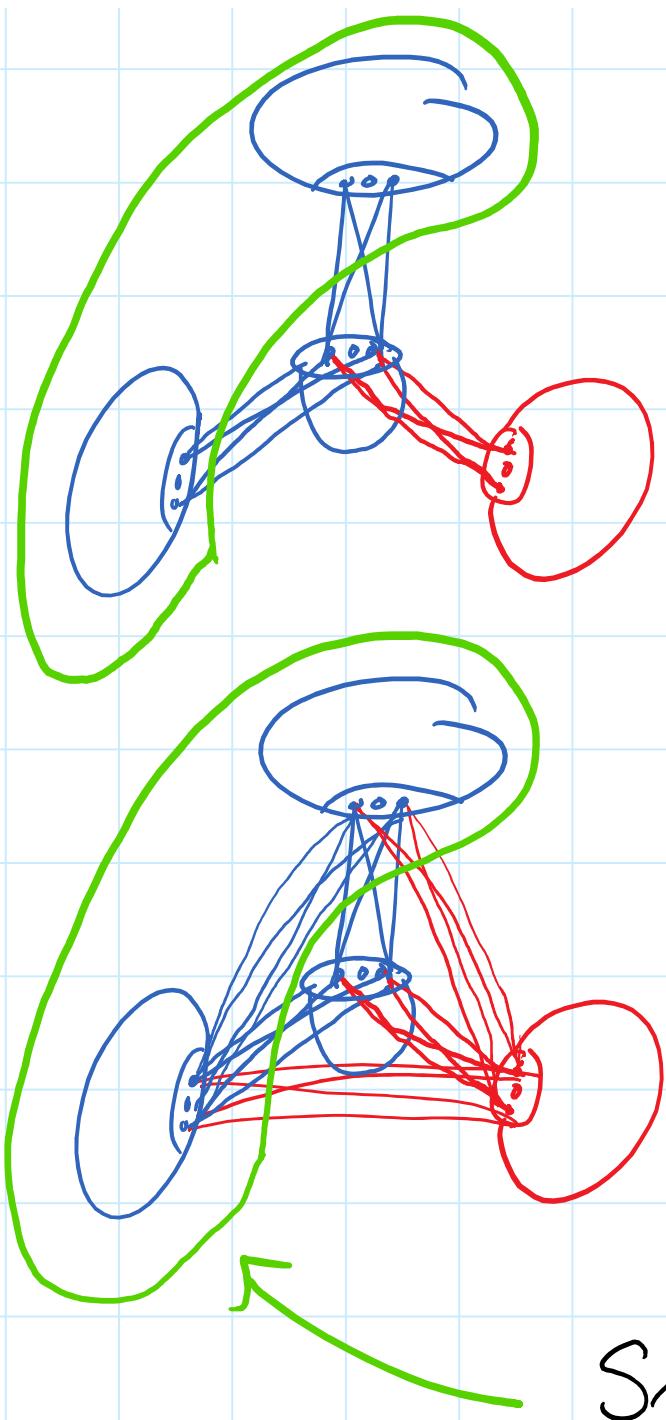
- As the chord diagram is connected
there is a path between 2 non-adjacent
vertices

- So one gets (by minimality):









Splits of G

